

# Hopfield Methods: Application to Optimization of Distributed Energy Resources

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Santa Fe, New Mexico USA





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PEREZ



Dr. Hongcai  
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Dr. Milad  
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Laurel  
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Zach  
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Saehong  
PARK



Dong  
ZHANG



Victoria  
CHENG



Sangjae  
BAE



Bertrand  
TRAVACCA



Mathilde  
BADOUAL



Zhe  
ZHOU



Tianyu  
HU



Tianyu  
YANG



Ramon  
CRESPO



Armando  
DOMINGOS



Soomin  
WOO



Yiqi  
ZHAO



Dylan  
KATO



Teng  
ZENG



Patrick  
KEYANTUO



Aaron  
KANDEL



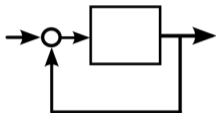
Pierre-François  
MASSIANI



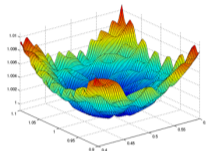
Emily  
YOU



FUNDAMENTAL  
RESEARCH



Dynamic Systems  
& Control

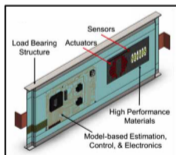


Optimization



Data Science

APPLICATIONS



Battery Management  
Systems (#BATT)



Automated, Connected, &  
Electric Vehicles (#ACES)



Distributed Energy  
Resources (#DER)



Laurel  
DUNN

Probabilistic Model for Assessing  
Distribution Grid Performance  
under Hazard Scenarios



Mathilde  
BADOUAL

Optimal Bidding Strategy on the  
Energy Market - A Reinforcement  
Learning Approach



Victoria  
CHENG

Virtual Inertia Frequency  
Regulation for Renewable  
Integration



Teng  
ZENG

Solving Overstay in PEV Charging  
Station Planning via Chance  
Constrained Optimization



Armando  
DOMINGOS

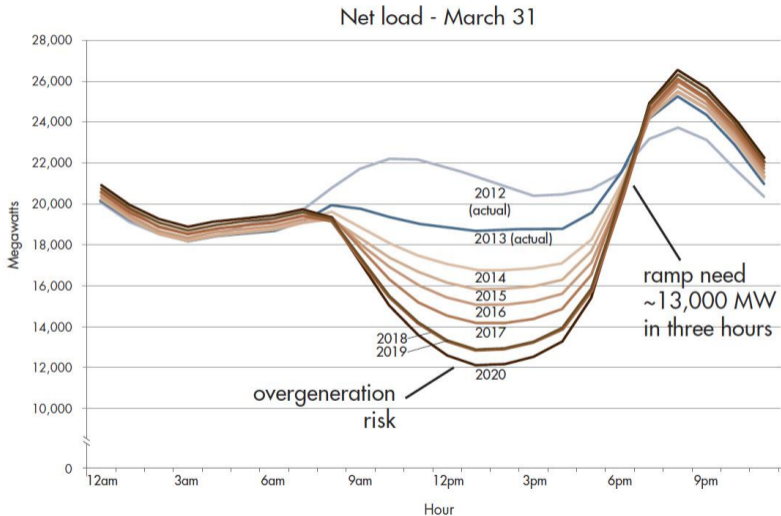
Multi-Armed Bandits  
DER Control



Bertrand  
TRAVACCA

THIS TALK

## The duck curve shows steep ramping needs and overgeneration risk



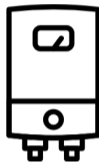
# Aggregate Flexible Loads into Virtual Power Plant



**Air  
Conditioning**



**Heater**



**Water  
Heater**



**Energy  
Storage  
System**



**Plug-in  
Electric  
Vehicles**

- Chen, Hashmi, Mathias, Busic, Meyn (2018)

# ★ PEV Charge Schedule Optimization is a MIP!

## eMotorWerks' Juicebox



Control  $\in$  *on* (40 A)  
or *off* (0 A)

## UC Berkeley Smart EV Charger *Richmond Field Station*



Control  $\in \{0 \text{ A}\} \cup [12 \text{ A}, 30 \text{ A}]$

Only  $\approx 10 \%$  of papers on large-scale optimization of PEVs model DISCRETE charging rates!

# ★ Problem Statement

Consider a mixed integer nonlinear program (MINLP):

$$\text{minimize} \quad f(\mathbf{x}) \quad (1)$$

$$\text{subject to:} \quad g_j(\mathbf{x}) \leq 0, \quad j = 1, \dots, m \quad (2)$$

$$\mathbf{x}_i \in \{0, 1\}, \quad i = 1, \dots, p < n \quad (3)$$

$$0 \leq \mathbf{x}_i \leq 1, \quad i = p + 1, \dots, n \quad (4)$$

$\mathbf{x} \in \mathbb{R}^n$  is the optimization variable

the first  $p < n$  variables must be binary

$f(\cdot) : \mathbb{R}^n \rightarrow \mathbb{R}$  is quadratic and  $L_f$  – smooth

$g_j(\cdot) : \mathbb{R}^n \rightarrow \mathbb{R}$  are quadratic and  $L_j$  – smooth

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## Challenge

Solve **LARGE-SCALE** MINLPs, e.g.  $n = 10^3, 10^4, 10^5, \dots$

P vs NP – Millenium Prize Problem



CLAY  
MATHEMATICS  
INSTITUTE

# Existing Convex Relaxation Methods

- 1 Binary relaxation
- 2 Lagrangian relaxation
- 3 Semi-definite relaxation
- 4 McCormick relaxations  
[McCormick '76][Nagarajan '16]
- 5 SoA Branch-and-Bound  
(linear relaxation) [Belotti '08]
- 6 SoA Branch-and-Cut  
[Achterberg '08]
- 7 Quadratic Convex  
relaxations [Hijazi '17]
- 8 Polyhedral relaxations for  
MIMFs [Nagarajan '18]

# Existing Convex Relaxation Methods

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- 2 Lagrangian relaxation
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MIMFs [Nagarajan '18]

## Stochastic approach to recover integer constraint:

Let  $x^r$  be sol'n to binary relaxation. Feasible  $x$  can be drawn randomly from  $\{0, 1\}$  following Bernoulli distribution  $\mathcal{B}(x^r)$ .

This can be sub-optimal.

### Example

$$\text{minimize}_{x \in \{0,1\}} \left(x - \frac{1}{4}\right)^2 = \frac{1}{16} \quad (x^* = 0 \text{ is opt. sol'n})$$

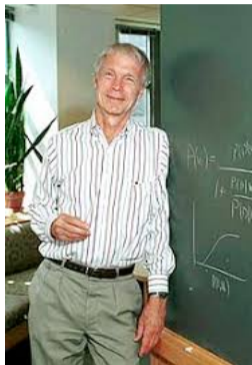
If we apply binary relaxation, we get  $x^r = \frac{1}{4}$  and  
 $\mathbb{E}_{x \sim \mathcal{B}(x^r)} \left(x - \frac{1}{4}\right)^2 = \frac{3}{16} > \frac{1}{16} !$

# Outline

- 1 Hopfield Methods - What are they?
- 2 Theoretical Analysis
- 3 Dual Hopfield Method
- 4 Example and Application

# A short history of Hopfield Networks

- (1982) J. J. Hopfield used neural nets to model collaborative computations
- (1985) J. J. Hopfield showed that neural nets can be used to solve optimization problems
- (1990's – 2000's) Hopfield methods became very popular for solving MIQPs in power systems optimization
- In literature, power system researchers admit they didn't fully understand *why* Hopfield methods work well.



# The Hopfield Method

Consider MINLP

$$\text{minimize} \quad f(\mathbf{x}) \quad (5)$$

$$\text{subject to:} \quad \mathbf{x}_i \in \{0, 1\}, \quad i = 1, \dots, p < n \quad (6)$$

$$0 \leq \mathbf{x}_i \leq 1, \quad i = p + 1, \dots, n \quad (7)$$

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Hopfield method follows dynamics:

$$\frac{d}{dt} \mathbf{x}_H(t) = -\nabla f(\mathbf{x}(t)); \quad \mathbf{x}_H(0) = \mathbf{x}(0) \in (0, 1)^n \quad (8)$$

$$\mathbf{x}(t) = \sigma(\mathbf{x}_H(t)) \quad (9)$$

where  $\sigma(\cdot) : \mathbb{R}^n \rightarrow [0, 1]^n$  is an “activation function” defined element-wise as:

$$\sigma(\mathbf{x}) : \mathbf{x} \mapsto [\sigma_1(x_1), \dots, \sigma_n(x_n)]$$

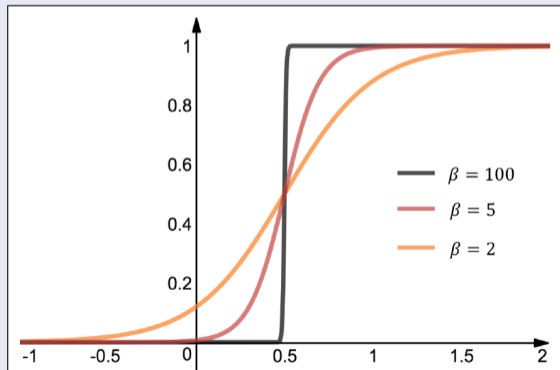
# What is activation function $\sigma(x)$ ?

- strictly increasing
- $\sigma(\cdot) \in \mathbb{C}^1$  with Lipschitz constant  $L_{\sigma_i}$

## Example: tanh

$$\sigma_i(x) = \frac{1}{2} \tanh(\beta_i(x - \frac{1}{2})) + \frac{1}{2}; \quad \beta_i > 0$$

“soft projection operator” from  $\mathbb{R}$  to  $\{0, 1\}$



## ★ ★ Hopfield Method $\Rightarrow$ Nonlinear Gradient Flow

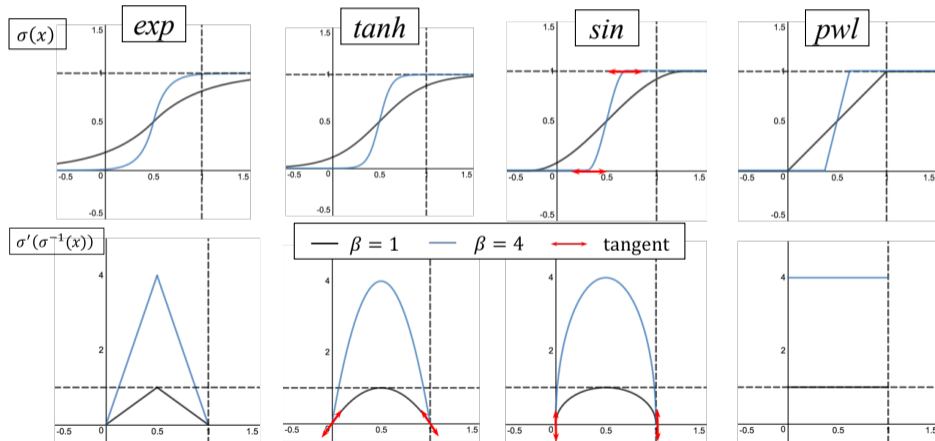
If  $\sigma(\cdot)$  is a homeomorphism, then a **nonlinear gradient flow emerges!**

$$\frac{d}{dt} \mathbf{x}(t) = -\sigma'(\sigma^{-1}(\mathbf{x}(t))) \odot \nabla f(\mathbf{x}(t)) \quad (10)$$

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# Discretize time dynamics

Forward Euler time discretization of Hopfield dynamics:

$$\mathbf{x}_H^{k+1} = \mathbf{x}_H^k - \alpha^k \nabla f(\mathbf{x}^k); \quad \mathbf{x}_H^0 = \mathbf{x}^0 \in (0, 1)^n \quad (11)$$

$$\mathbf{x}^k = \sigma(\mathbf{x}_H^k) \quad (12)$$

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For quadratic  $f(\mathbf{x}) = \frac{1}{2} \mathbf{x}^T \mathbf{Q} \mathbf{x}$

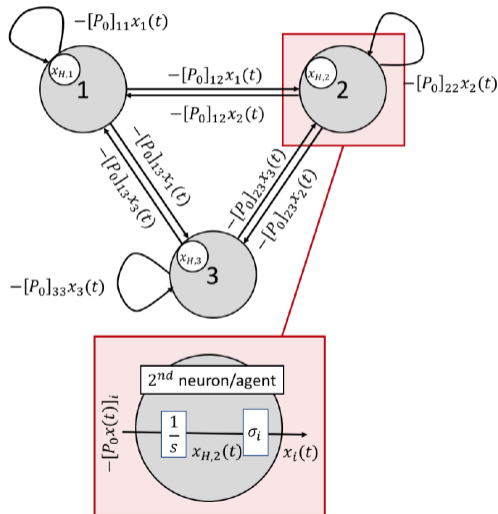
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$$\mathbf{x}^k = \sigma(\mathbf{x}_H^k) \quad (14)$$

# Graphical Interpretation of Hopfield Method

## Forward Simulation of Hopfield Neural Net!

- Undirected weighted graph
- $n$  nodes, one for each  $x_i$
- Each node has internal ( $x_{H,i} \in \mathbb{R}$ ) and external ( $x_i \in \mathbb{R}$ ) states
- Weights  $[P_0]_{ij}$  are elements of gradients of obj fcn



# ★ Hopfield vs Projected Gradient Descent

## Hopfield

$$x_H^{k+1} = x_H^k - \alpha^k \nabla f(x^k) \quad (15)$$

$$x^k = \sigma(x_H^k) \quad (16)$$

## Projected Gradient Descent

$$x_H^{k+1} = x^k - \alpha^k \nabla f(x^k) \quad (17)$$

$$x^k = \text{Proj}_{[0,1]}(x_H^k) \quad (18)$$

# ★ Hopfield vs Projected Gradient Descent

## Hopfield

$$x_H^{k+1} = x_H^k - \alpha^k \nabla f(x^k) \quad (15)$$

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## Projected Gradient Descent

$$x^{k+1} = \text{Proj}_{[0,1]}(x^k - \alpha^k \nabla f(x^k)) \quad (17)$$

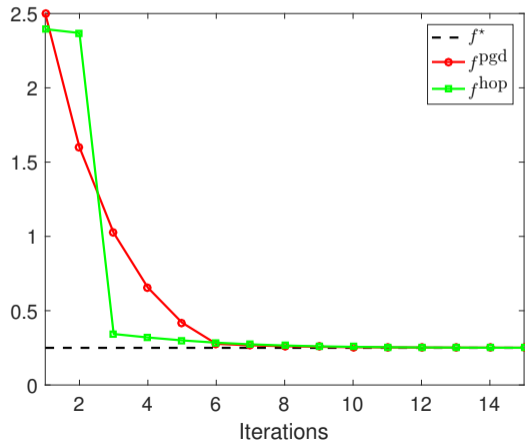
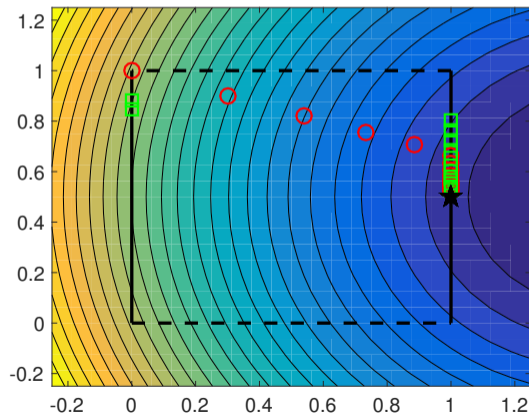
No dynamics!

# Simple Comparison

$$\text{minimize}_{x_1, x_2} \quad (x_1 - 1.5)^2 + (x_2 - 0.5)^2 \quad (18)$$

$$\text{subject to:} \quad x_1 \in \{0, 1\} \quad (19)$$

$$0 \leq x_2 \leq 1 \quad (20)$$



# Outline

- 1 Hopfield Methods - What are they?
- 2 Theoretical Analysis**
- 3 Dual Hopfield Method
- 4 Example and Application

# Continuous Improvement to a Fixed Point

## Theorem 1: Continuous Improvement

The Hopfield method yields monotonically decreasing iterates,  $f(x^{k+1}) \leq f(x^k)$ ,  $\forall k$  if ...

- activation fcn has Lipschitz continuous first derivative:  $\sigma(\cdot) \in \mathbb{C}^1$  (exp, tanh, sin, pwl)
- step-size  $\alpha^k$  follows an appropriately decreasing schedule

Specifically, the incremental improvement is bounded by:

$$0 \leq f(x^k) - f(x^{k+1}) \leq 0.5\alpha^k \cdot \nabla f(x^k)^T \Sigma^k \nabla f(x^k) \quad \text{where } \Sigma^k = \text{diag}(\sigma'(x_H^k))$$

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## Corollary: Convergence within a set

There exists a  $f^\dagger$  such that  $f(x^k) \rightarrow f^\dagger$  as  $k \rightarrow \infty$ , and  $x^k$  converges to the (non-empty) set

$$\mathcal{X} = \left\{ x \in [0, 1]^n \mid x_i \in \{0, 1\} \quad \text{OR} \quad \frac{\partial}{\partial x_i} f(x) = 0, \quad i = 1, \dots, p \right\} \quad (21)$$

### Remarks:

- Set  $\mathcal{X}$  includes true minimizer  $x^*$ , but  $x^k \rightarrow x^*$  not guaranteed

## Theorem 2: Sub-linear convergence

If  $f(x)$  is convex and  $\sigma(\cdot)$  is smooth and verifies

$$\sigma'(\sigma^{-1}(x)) \geq \min \{|x|, |1 - x|\}, \quad x \in [0, 1]^n \quad (22)$$

then,

- $f(x^k) - f^\dagger = \mathcal{O}\left(\frac{1}{k^r}\right)$ , with  $0 < r < 1$
- To achieve precision  $\varepsilon$ , the worst case number of iterations is  $2Mn/(\beta^2\varepsilon)$ 
  - $M$  is upper-bound on Hessian:  $\nabla^2 f(x) \preceq MI$
  - $n$  is number of variables  $x \in \mathbb{R}^n$
  - $\beta$  is “hardness” of activation function

**Remark:** Slower than gradient descent, for which convergence is guaranteed at rate  $\mathcal{O}\left(\frac{1}{k}\right)$

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# Dual Hopfield Method

So far, we have considered Hopfield methods to approximately solve

$$\text{minimize} \quad f(\mathbf{x}) \quad (23)$$

$$\text{subject to:} \quad 0 \leq x_i \leq 1 \quad i = 1, \dots, n \quad (24)$$

$$x_i \in \{0, 1\} \quad i = 1, \dots, p < n \quad (25)$$

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$$x_i \in \{0, 1\} \quad i = 1, \dots, p < n \quad (25)$$

We now consider **inequality constraints**:

$$\text{minimize} \quad f(x) \quad (26)$$

$$\text{subject to:} \quad g_j(x) \leq 0, \quad j = 1, \dots, m \quad (27)$$

$$0 \leq x_i \leq 1 \quad i = 1, \dots, n \quad (28)$$

$$x_i \in \{0, 1\} \quad i = 1, \dots, p < n \quad (29)$$

# Dual Hopfield Method

Apply Lagrangian relaxation

**Idea:** Instead of considering the “full” Lagrangian relaxation, consider

$$L(\mathbf{x}, \boldsymbol{\mu}) = f(\mathbf{x}) + \sum_{j=1}^m \mu_j g_j(\mathbf{x}) \quad (30)$$

# Dual Hopfield Method

Apply Lagrangian relaxation

**Idea:** Instead of considering the “full” Lagrangian relaxation, consider

$$L(\mathbf{x}, \boldsymbol{\mu}) = f(\mathbf{x}) + \sum_{j=1}^m \mu_j g_j(\mathbf{x}) \quad (30)$$

Then the *dual function* is

$$D(\boldsymbol{\mu}) = \min_{\mathbf{x}} \quad L(\mathbf{x}, \boldsymbol{\mu}) = f(\mathbf{x}) + \sum_{j=1}^m \mu_j g_j(\mathbf{x}) \quad (31)$$

$$\text{subject to:} \quad 0 \leq x_i \leq 1 \quad i = 1, \dots, n \quad (32)$$

$$x_i \in \{0, 1\} \quad i = 1, \dots, p < n \quad (33)$$

which is amenable to Hopfield method, given  $\boldsymbol{\mu}$ .

# Dual Ascent via Hopfield

Then solve the Dual Problem:

$$\max_{\mu \geq 0} D(\mu) \quad (34)$$

$$D(\mu) = \min_x L(x, \mu) = \min_x f(x) + \sum_{j=1}^m \mu_j g_j(x) \quad (35)$$

# Dual Ascent via Hopfield

Then solve the Dual Problem:

$$\max_{\mu \geq 0} D(\mu) \tag{34}$$

$$D(\mu) = \min_x L(x, \mu) = \min_x f(x) + \sum_{j=1}^m \mu_j g_j(x) \tag{35}$$

Run Hopfield method to approximately solve  $D(\mu) = \min_x L(x, \mu)$ .

Suppose  $x^*(\mu) = \arg \min_x L(x, \mu)$ .

The subgradient of  $D(\mu)$  along dimension  $j$ :  $g_j(x^*(\mu)) \in \partial_j D(\mu)$

# Dual Hopfield Method

## The Algorithm

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**Algorithm 1** Dual (sub)-gradient Ascent via Hopfield Method

---

*Initialize*  $\lambda^0 \geq 0$ ; Choose  $\beta > 0$

**for**  $k = 0, 1, \dots, k_{\max}$

**(1)** use Hopfield method to approximately compute dual function

**for**  $\ell = 0, \dots, \ell_{\max}$

$$x_H^{\ell+1} = x_H^\ell - \alpha^\ell \nabla_x L(x^\ell, \mu^k)$$

$$x^\ell = \sigma(x_H^{\ell+1})$$

$$x_{\text{hop}}^k \leftarrow x^\ell$$

**until** stopping criterion is met

**(2)** update dual variable  $\mu$  via (sub)-gradient ascent

$$\mu^{k+1} = \mu^k + \beta^k \sum_{j=1}^m g_j(x_{\text{hop}}^k(\mu^k))$$

**end for**

---

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# Examples: Random MIQPs

Consider solving MIQP w.r.t.  $x \in \mathbb{R}^n$

$$\text{minimize} \quad \frac{1}{2}x^T Qx + R^T x \quad (36)$$

$$\text{subject to:} \quad Ax \leq b \quad (37)$$

$$A_{eq}x = b_{eq} \quad (38)$$

$$lb \leq x \leq ub \quad (39)$$

$$x_i \in \{0, 1\}, \quad i = 1, \dots, p \quad (40)$$

- Randomly generated parameters  $Q, R, A, b, A_{eq}, b_{eq}, lb, ub$  for each  $n$
- Number of constraints also randomized

# Comparative Analysis

All problems solved on Matlab:

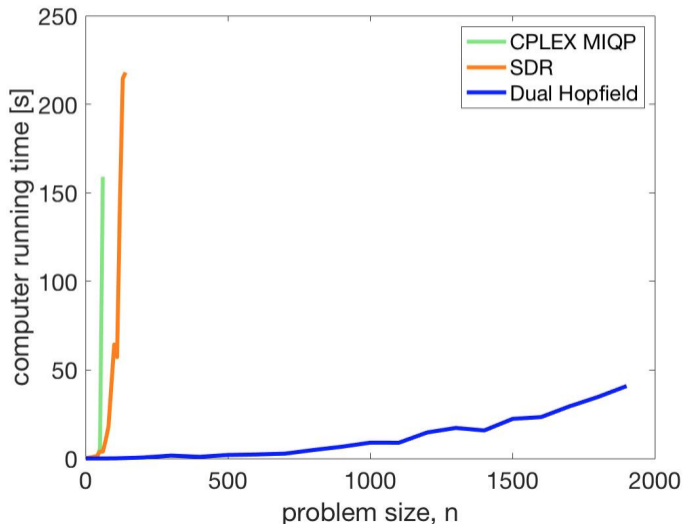
- **CPLEX MIQP**: using function *cplexmiqp* developed by IBM
- Binary Relaxation via **CPLEX QP** : using function *cplexqp*
- Semi-definite relaxation (**SDR**): corresponding SDP solved using CVX
- **Hopfield**: Dual Ascent Hopfield Method uses dual variables from *cplexqp*

For each method, we compute:

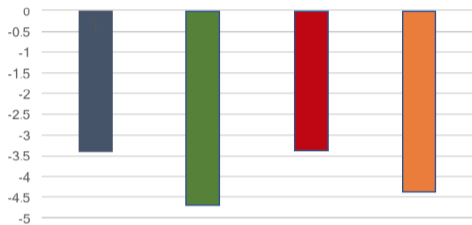
- computer running time [sec]
- constraint violations (CV):
  - binary CV:  $\frac{1}{p} \sum_{i=1}^p d(x_i, \{0, 1\})$
  - inequality CV:  $\frac{1}{m} \sum_{j=1}^m |[Ax - b]_j|$
  - equality CV:  $\frac{1}{\ell} \sum_{k=1}^{\ell} |[A_{eq}x - b_{eq}]_k|$
- objective function value

# Comparative Analysis

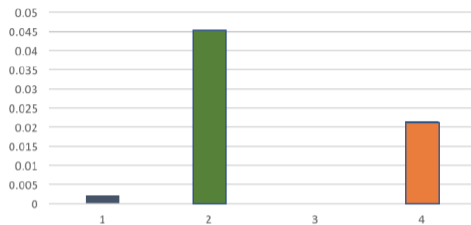
Computer running time



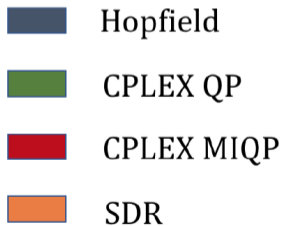
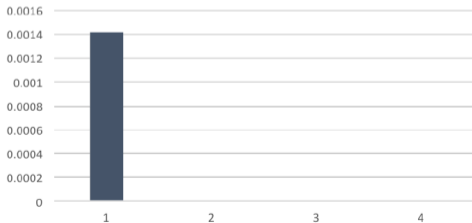
### Objective value



### Binary CV



### equality CV



# Application: Optimal Economic Dispatch of DERs

Consider  $n$  generators with cost:  $f_i(x_i) = c_i x_i^2 + b_i x_i + a_i$ , with  $a_i, b_i, c_i \geq 0$ .

The first  $p$  generators can only make binary decisions. That is:

- $\forall i \in \{1, p\}$  we have  $x_i \in \{P_{i,\min}, P_{i,\max}\}$
- $\forall i \in \{p+1, n\}$  we have  $x_i \in [P_{i,\min}, P_{i,\max}]$

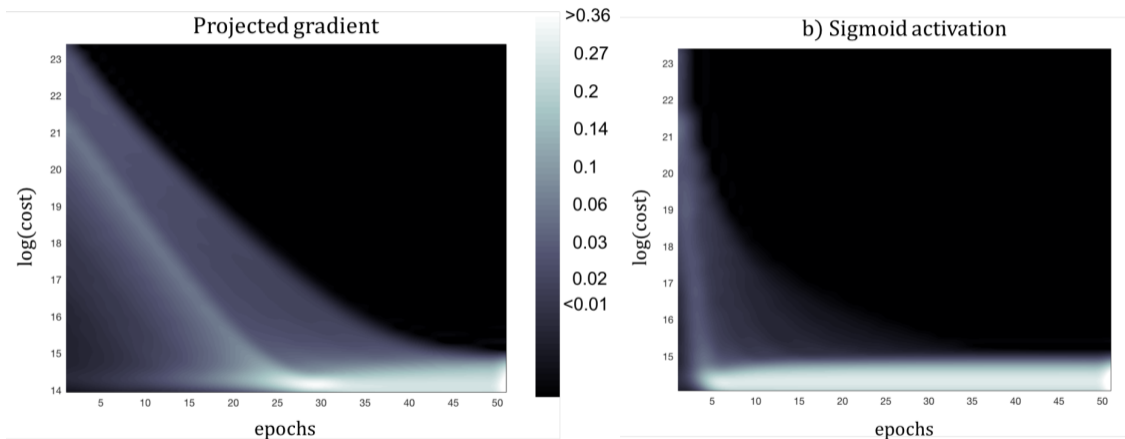
## Problem Statement

Find the optimal dispatch for generators to minimize cost and meet demand:

$$\begin{aligned} &\text{minimize} && \sum_i^n f_i(x_i) \\ &\text{subject to:} && \sum_i^n x_i = D \\ &&& \text{(constraints above)} \end{aligned}$$

Simulation parameters:  $n = 1000$  generators. Other parameters randomly generated. We perform 5000 Monte-Carlo simulations.

# Monte Caroline Simulation Results



## SUMMARY

- Hopfield Methods for large-scale MINLPs – *An old heuristic with new analysis!*

## EXTENSIONS

- Alternative descent direction
- Nesterov acceleration
- Chance constraints
- Distributed algorithms via dual decomposition
- ...

## ON-GOING / FUTURE

- Application to Large-Scale PEV Charge Scheduling
- More comprehensive comparative analysis
- Open source codes! `hmip`

# VISIT US!

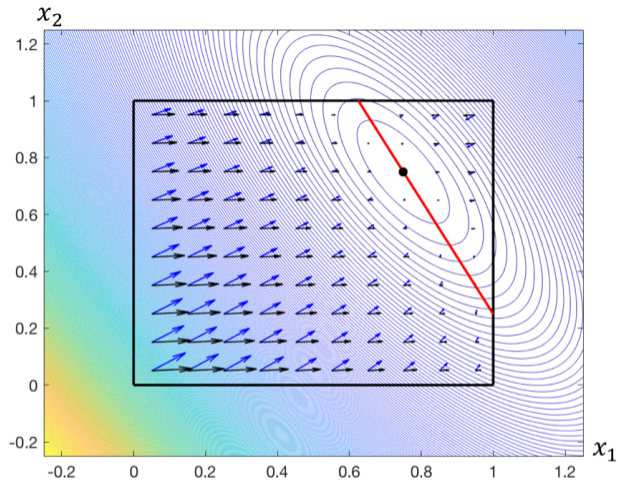
Energy, Controls, and Applications Lab (eCAL)

[ecal.berkeley.edu](http://ecal.berkeley.edu)

[smoura@berkeley.edu](mailto:smoura@berkeley.edu)



# APPENDIX SLIDES



$$\text{blue arrow} \rightarrow \nabla f_0(x) \quad \text{red line} \rightarrow \nabla f_0(x)_1 = 0$$

$$\text{black arrow} \rightarrow \sigma'(\sigma^{-1}(x)) \odot \nabla f_0(x) \quad \bullet \rightarrow \nabla f_0(x) = 0$$

# Existing Methods

## Convex Relaxation #1: Binary Relaxation

### **Stochastic approach to recover integer constraint:**

Let  $x^r$  be solution to binary relaxation. Feasible  $x$  can be drawn randomly from  $\{0, 1\}$  following Bernoulli distribution  $\mathcal{B}(x^r)$ .

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### Example

$$\text{minimize}_{x \in \{0,1\}} \left(x - \frac{1}{4}\right)^2 = \frac{1}{16} \quad (x^* = 0 \text{ is the optimal solution})$$

If we apply binary relaxation, we get  $x^r = \frac{1}{4}$  and  $\mathbb{E}_{x \sim \mathcal{B}(x^r)} \left(x - \frac{1}{4}\right)^2 = \frac{3}{16} > \frac{1}{16} !$

Other ideas:

- Branch & Bound, Branch & Cut

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## Convex Relaxation #2: Lagrangian Relaxation

Notice that  $x_i \in \{0, 1\}$  is equivalent to satisfying  $x_i(1 - x_i) = 0$

$$\text{minimize} \quad f(\mathbf{x}) \quad (41)$$

$$\text{subject to:} \quad g_j(\mathbf{x}) \leq 0, \quad j = 1, \dots, m \quad (42)$$

$$0 \leq \mathbf{x} \leq 1 \quad (43)$$

$$x_i(1 - x_i) = 0, \quad i = 1, \dots, p < n \quad (44)$$

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Form the *Lagrangian*:

$$L(\mathbf{x}, \mu, \underline{\mu}, \bar{\mu}, \lambda) = f(\mathbf{x}) + \sum_{j=1}^m \left[ \mu_j g_j(\mathbf{x}) + \underline{\mu}_j x_j + \bar{\mu}_j (1 - x_j) \right] + \sum_{i=1}^p \lambda_i x_i (1 - x_i) \quad (45)$$

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Define the (concave) *dual function* of  $\Lambda = [\mu, \underline{\mu}, \bar{\mu}, \lambda]$

$$D(\Lambda) = \min_{\mathbf{x} \in \mathbb{R}^n} L(\mathbf{x}, \mu, \underline{\mu}, \bar{\mu}, \lambda) \quad (46)$$

Weak duality approach: Solve convex program  $\max_{\Lambda} D(\Lambda)$

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$$\text{subject to:} \quad \frac{1}{2} \text{Tr}(Q_j X) + R_j^T x + S_j \leq 0, \quad j = 1, \dots, m \quad (48)$$

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$$X_{ii} = x_i, \quad i = 1, \dots, p < n \quad (50)$$

$$X = xx^T \quad (51)$$

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If  $Q, Q_i$  are positive semi-definite, then only  $X = xx^T$  makes this non-convex. Relax into convex inequality  $X \succeq xx^T$ . Using Schur complement:

$$X \succeq xx^T \Leftrightarrow \begin{bmatrix} X & x \\ x^T & 1 \end{bmatrix} \succeq 0 \quad (52)$$

This can be cast as a semi-definite program (SDP).