

Dynamical Structures in Iterative Decoding

Misha Stepanov

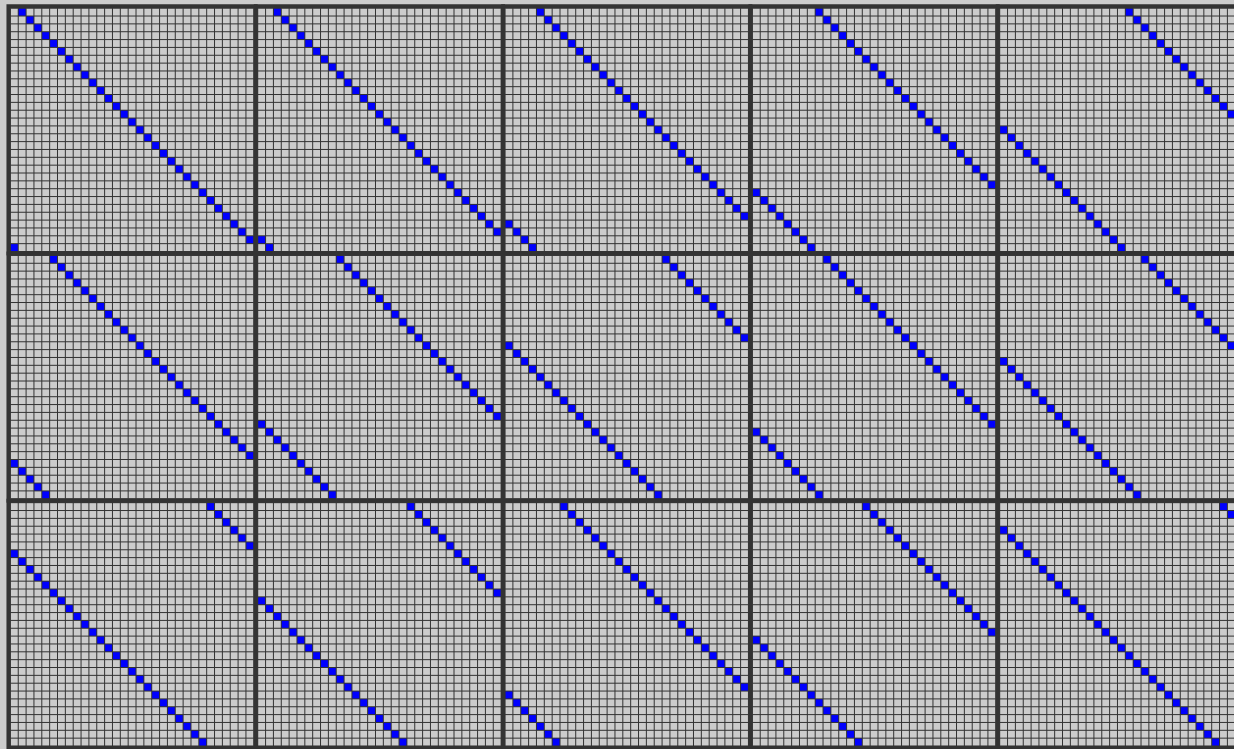
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Tanner's $[155, 64, 20]$ code

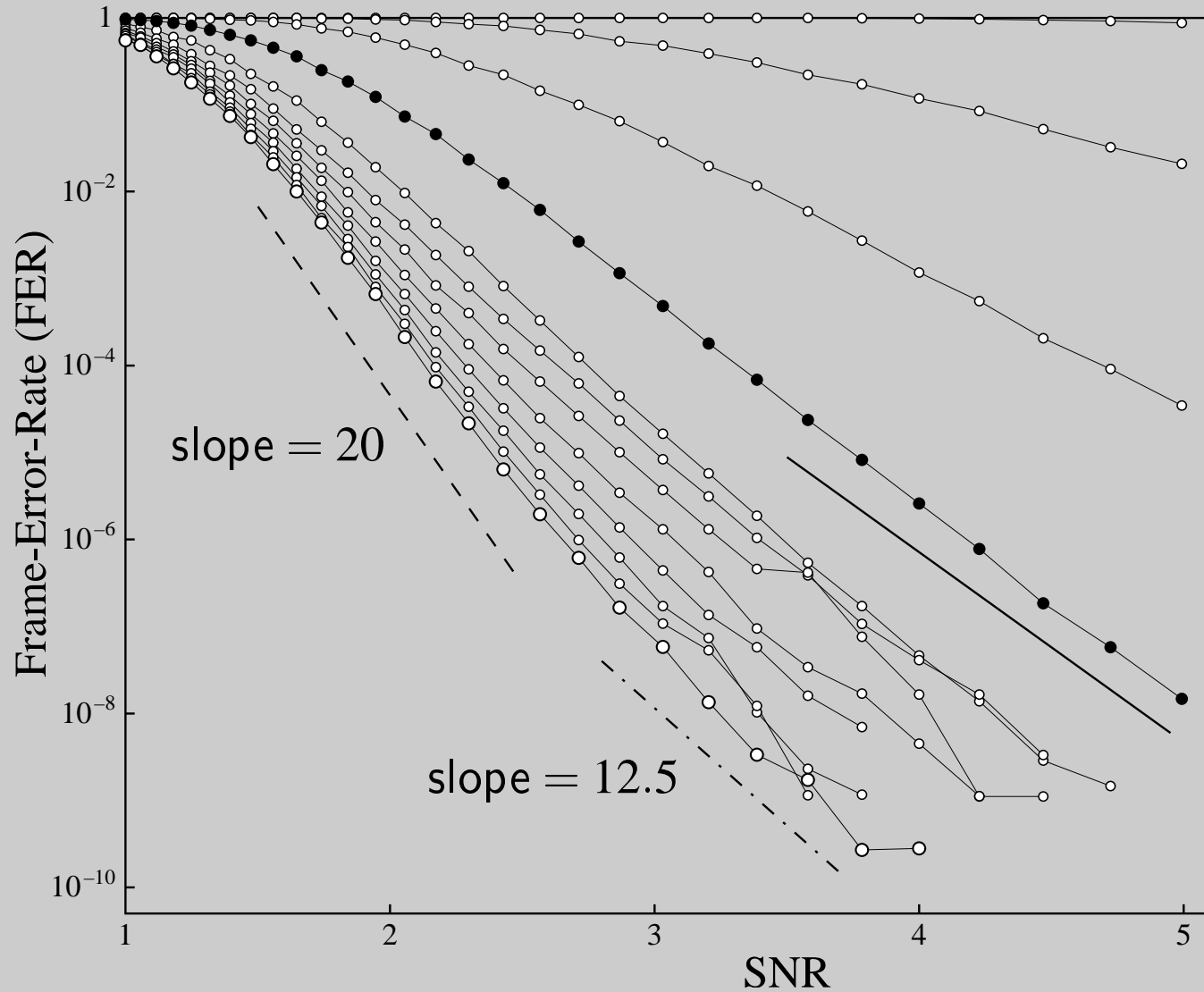
└── Hamming distance
└── informational bits
└── length of encoded message

Parity check matrix:



R.M. Tanner, D. Sridhara, T. Fuja, in *Proc. ISCTA 2001*
(Ambleside, UK, July 15–20, 2001), p. 365.

Frame-Error-Rate



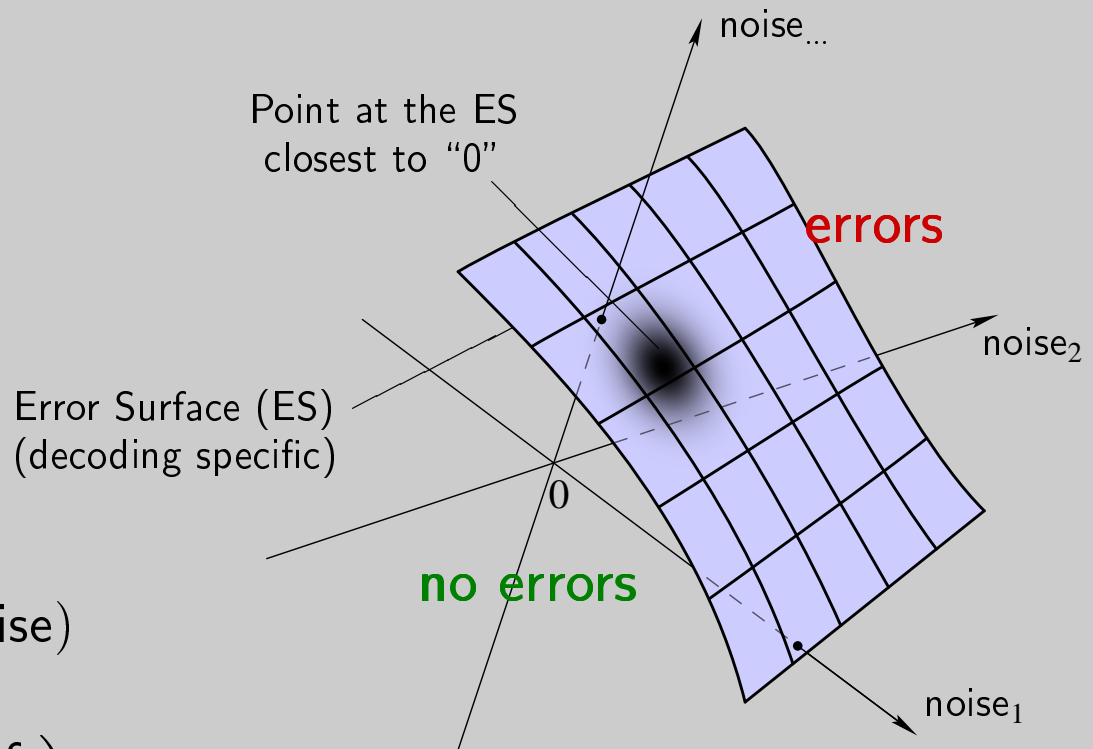
Instanton method

instanton method
saddle-point method
Laplace method
method of steepest descent
large deviations

$$\text{BER} = \int d(\text{noise}) \text{WEIGHT}(\text{noise})$$

$$\text{BER} \sim \text{WEIGHT} \left(\begin{array}{l} \text{optimal conf} \\ \text{of the noise} \end{array} \right)$$

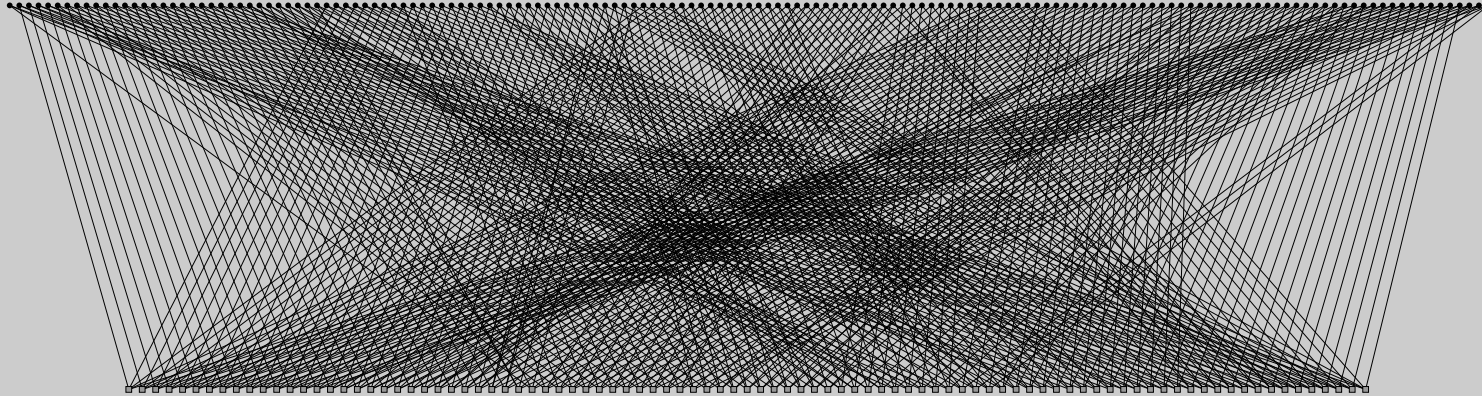
optimal conf
of the noise = Point at the ES
closest to "0"



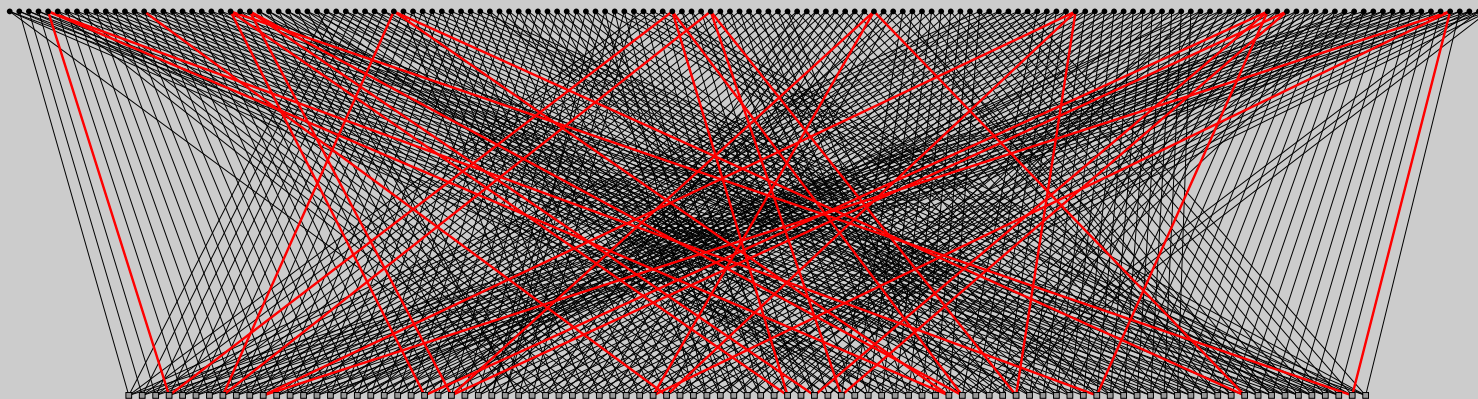
Chernyak, Chertkov, Stepanov, Vasic,
Phys. Rev. Lett. **93**, 198702 (2004)

Stepanov, Chertkov, Chernyak, Vasic,
Phys. Rev. Lett. **95**, 228701 (2005)
[cond-mat/0506037]

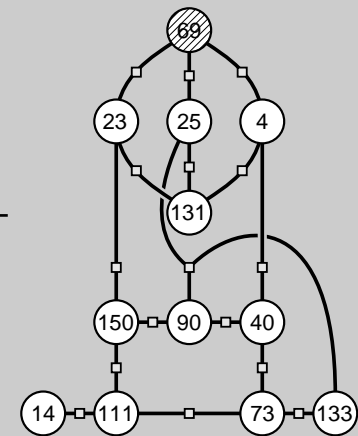
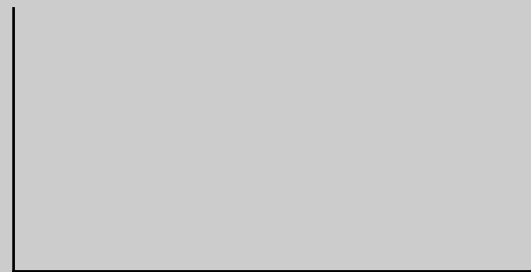
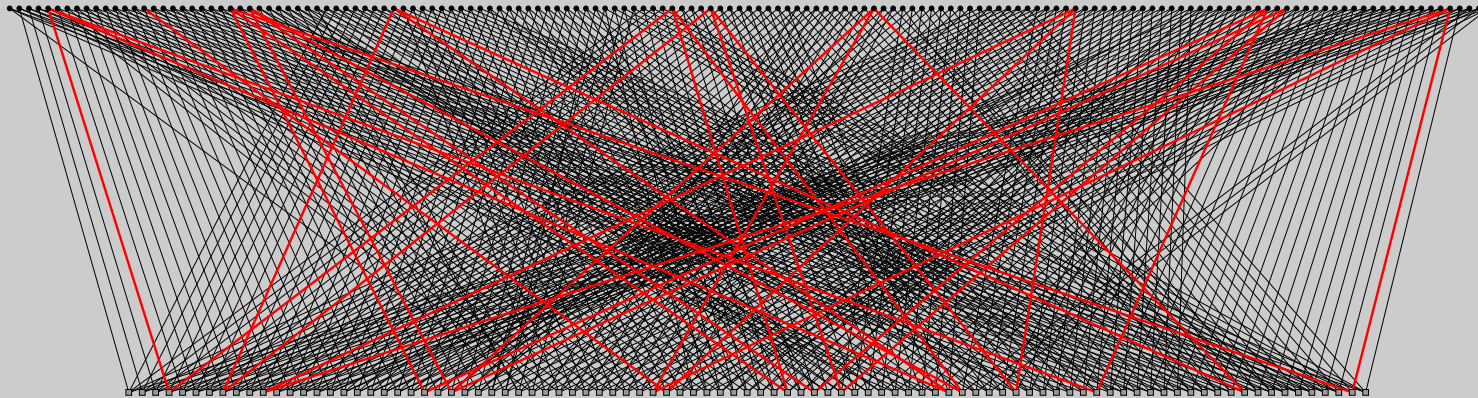
Tanner graph of $[155, 64, 20]$ code



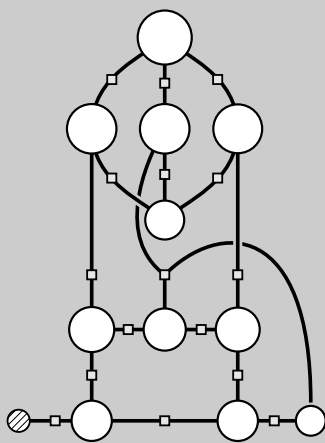
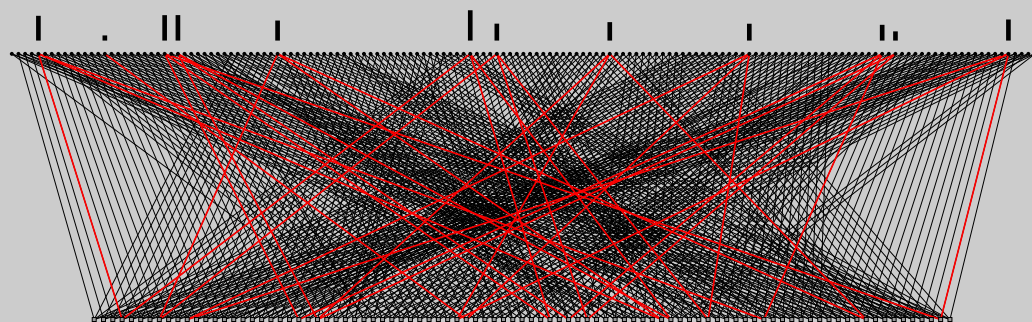
One special subgraph of Tanner graph



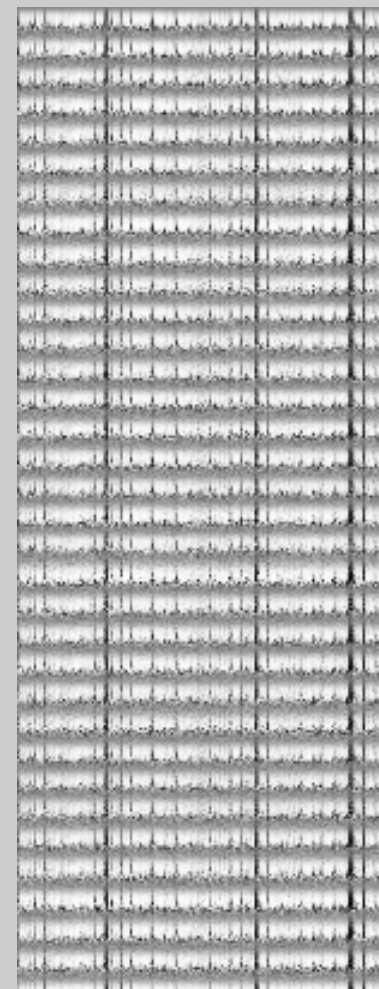
One special subgraph of Tanner graph



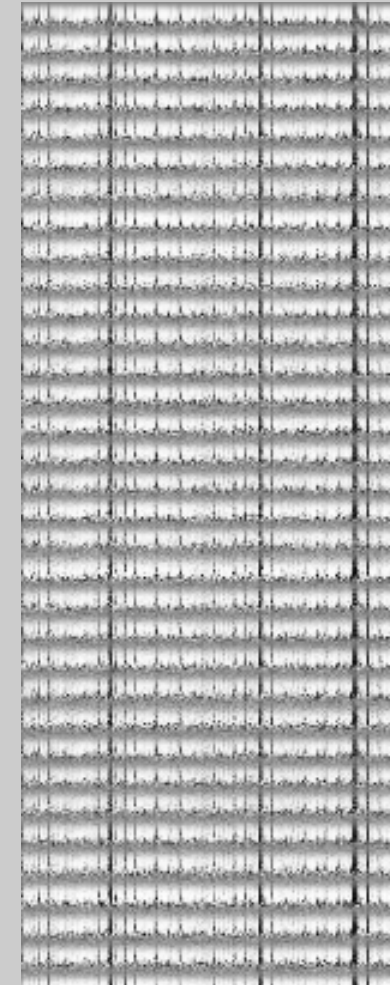
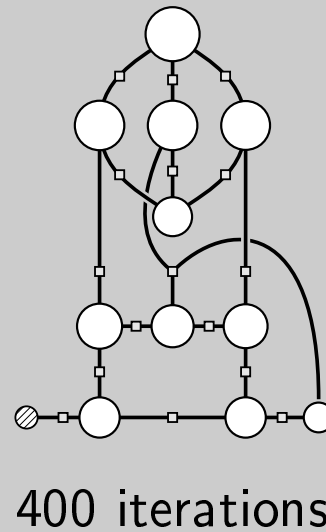
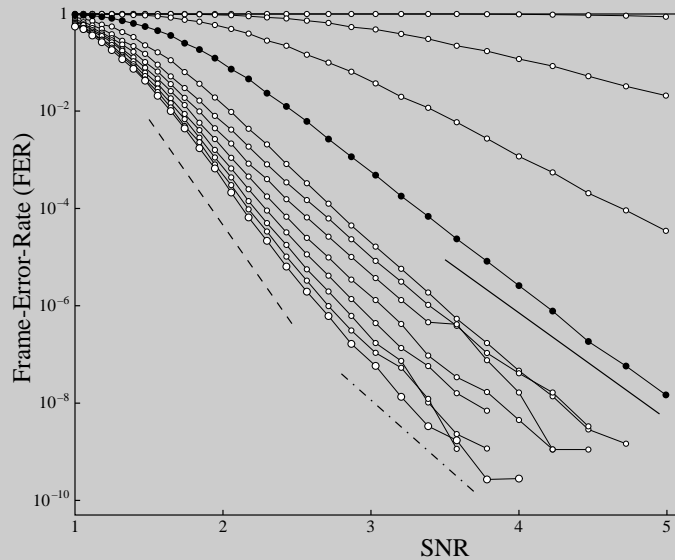
One special subgraph of Tanner graph



400 iterations



Instanton for Tanner's $[155, 64, 20]$ code



Effective distances:

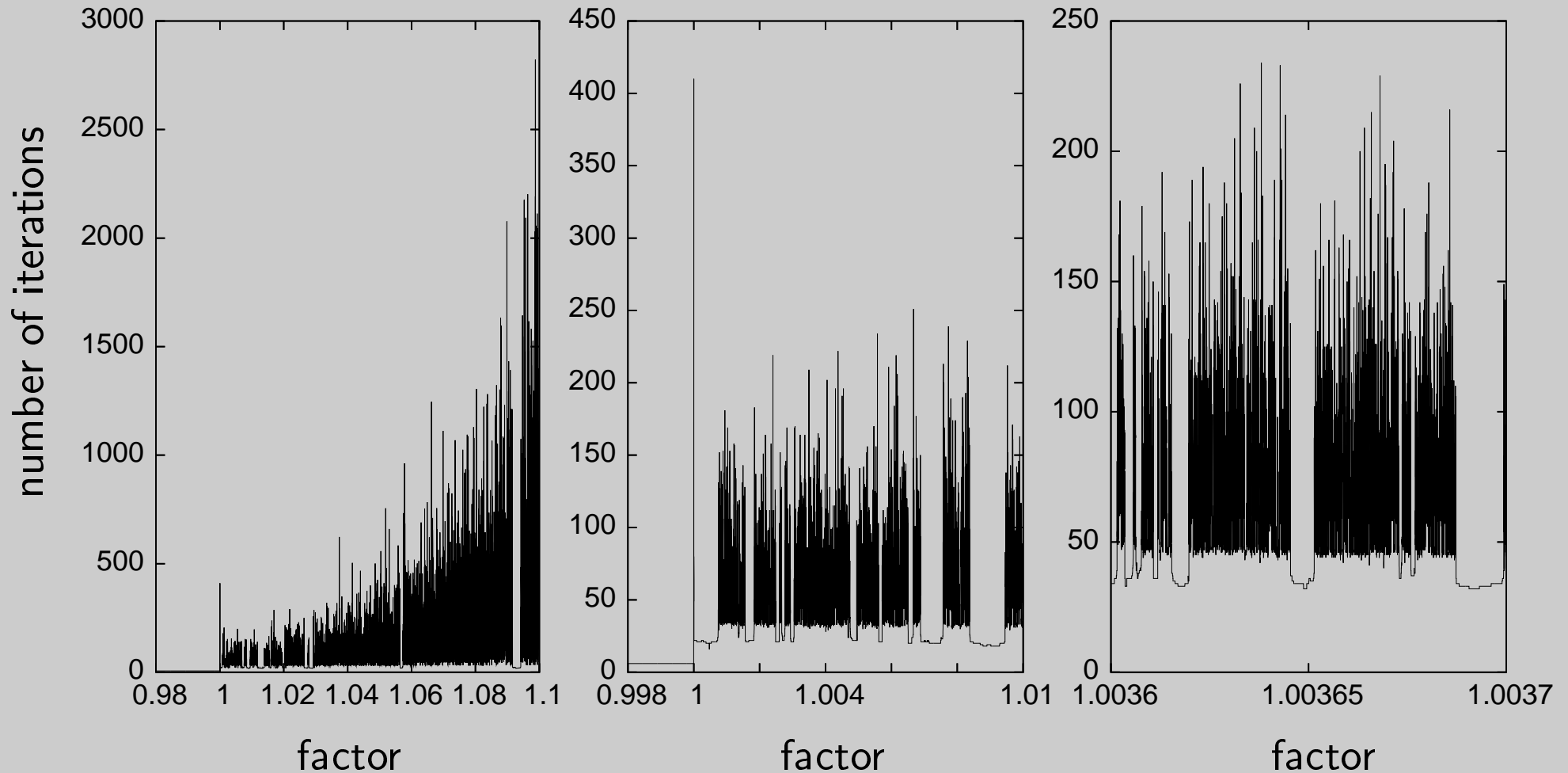
Iterative decoding: 12.5

Linear programming decoding: 16.4

Hamming distance: 20

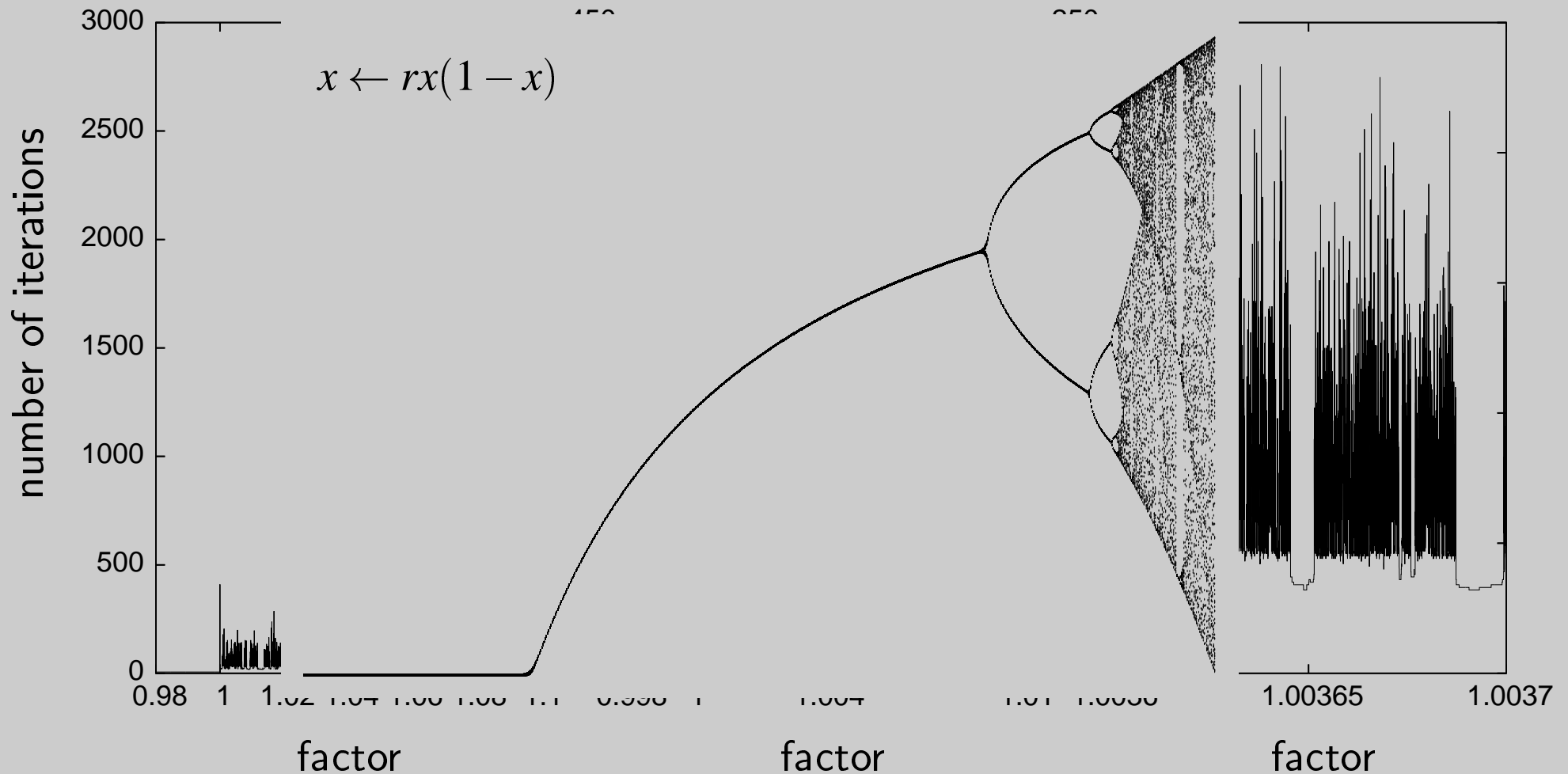
Instanton “robustness”

number of iterations until a successful decoding



Instanton “robustness”

number of iterations until a successful decoding



Smoothed (relaxed, damped) decoding

Iterative scheme (BP): $\eta_{i\alpha}^{(n+1)} = h_i + \sum_{\beta \ni i, \beta \neq \alpha} \tanh^{-1} \left(\prod_{j \in \beta} \tanh \eta_{j\beta}^{(n)} \right)$



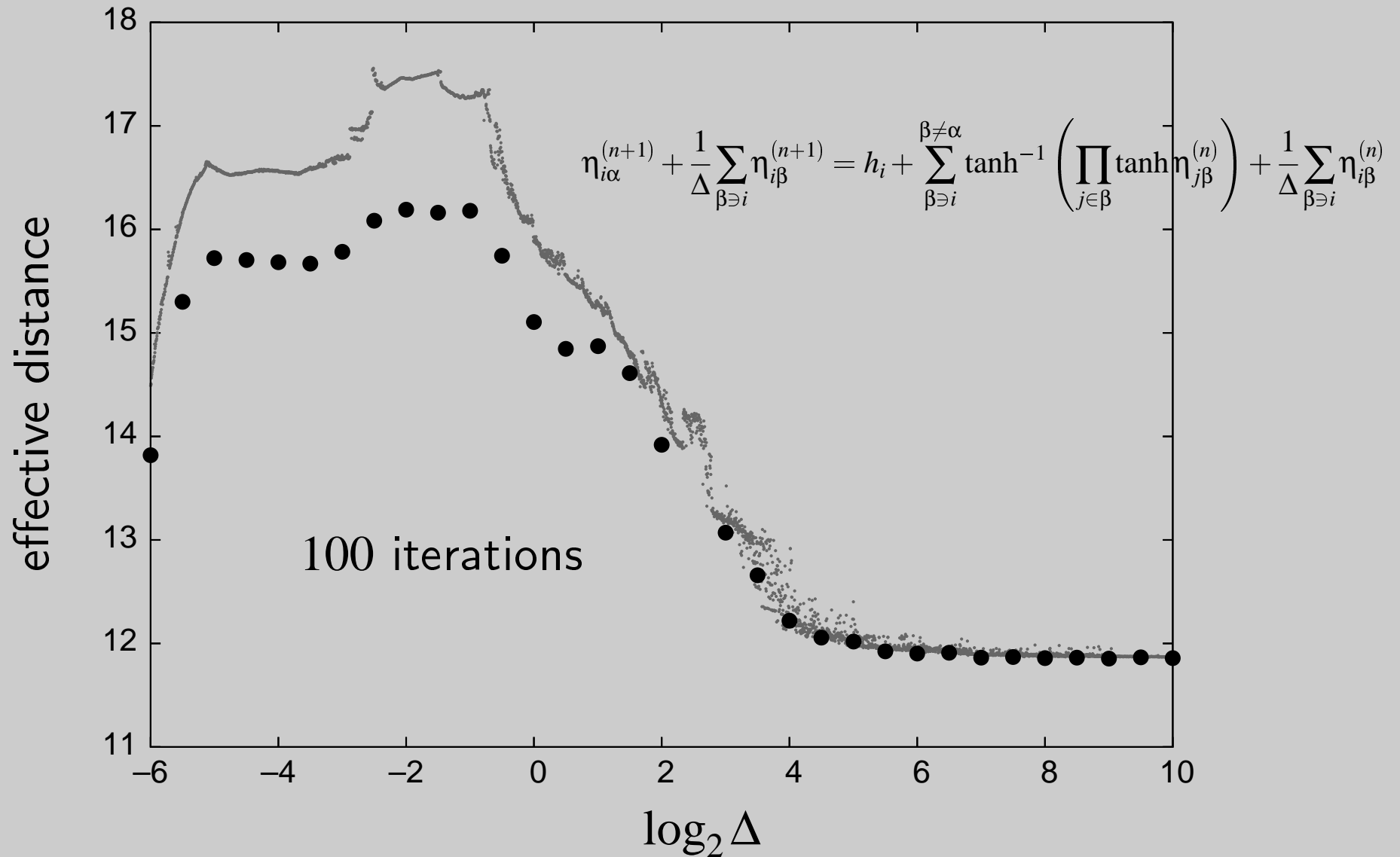
$$\eta_{i\alpha}^{(n+1)} + \frac{1}{\Delta} \sum_{\beta \ni i} \eta_{i\beta}^{(n+1)} = h_i + \sum_{\beta \ni i, \beta \neq \alpha} \tanh^{-1} \left(\prod_{j \in \beta} \tanh \eta_{j\beta}^{(n)} \right) + \frac{1}{\Delta} \sum_{\beta \ni i} \eta_{i\beta}^{(n)}$$

$\Delta \rightarrow \infty$ — standard BP

$\Delta \rightarrow 0$ — slow dynamics

Stepanov, Chertkov, Allerton 2006 [cs.IT/0607112]

Instantons effective distance

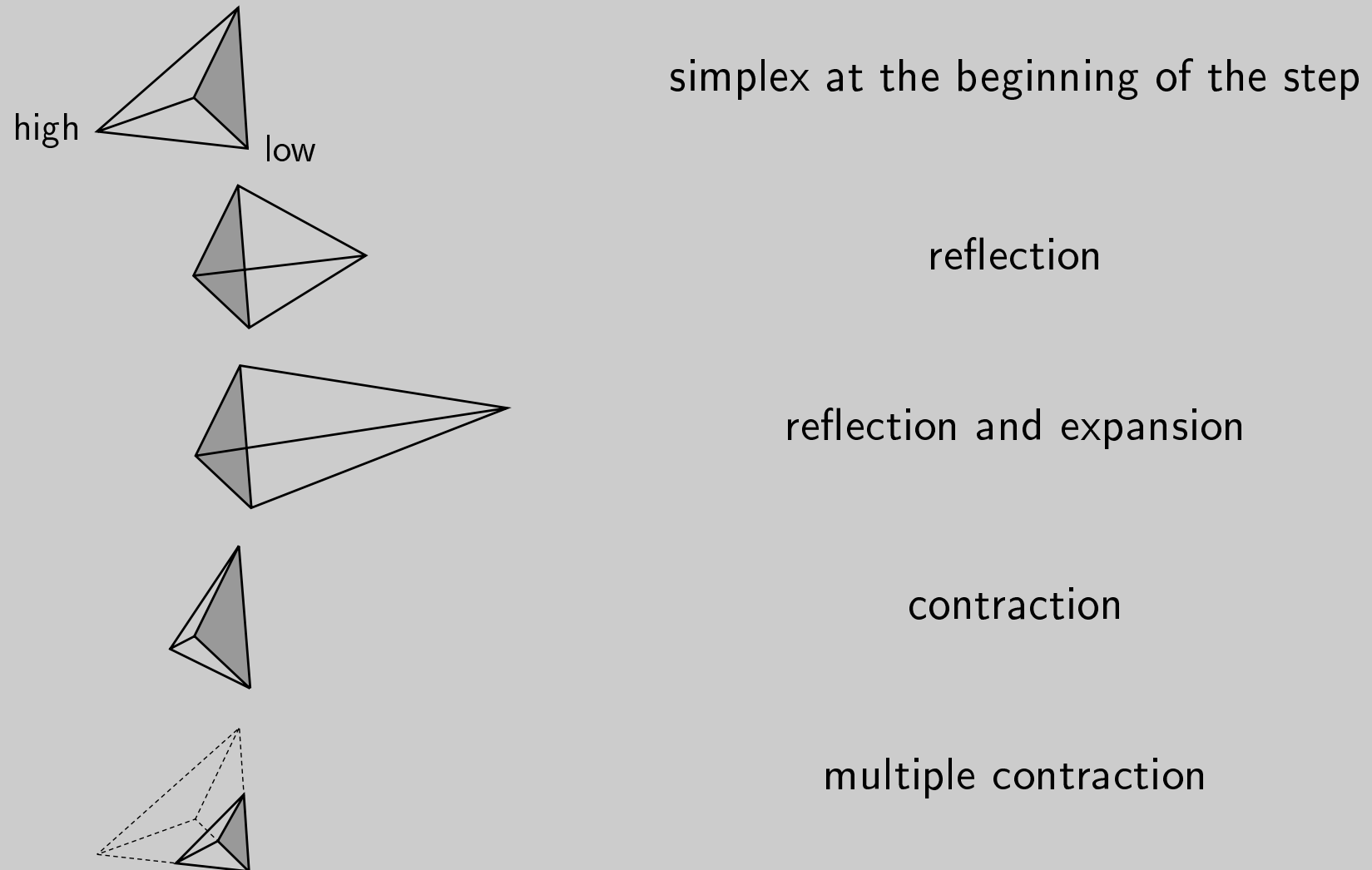


Summary

- the performance of iterative decoding is determined by most dangerous noise configurations (instantons)
- the fixed point of iterations in decoding is unstable, if the noise configuration is damaging
- the iterative decoding cycles on instantons
- making the iterations smoother helps (shifts the instantons to larger distances)

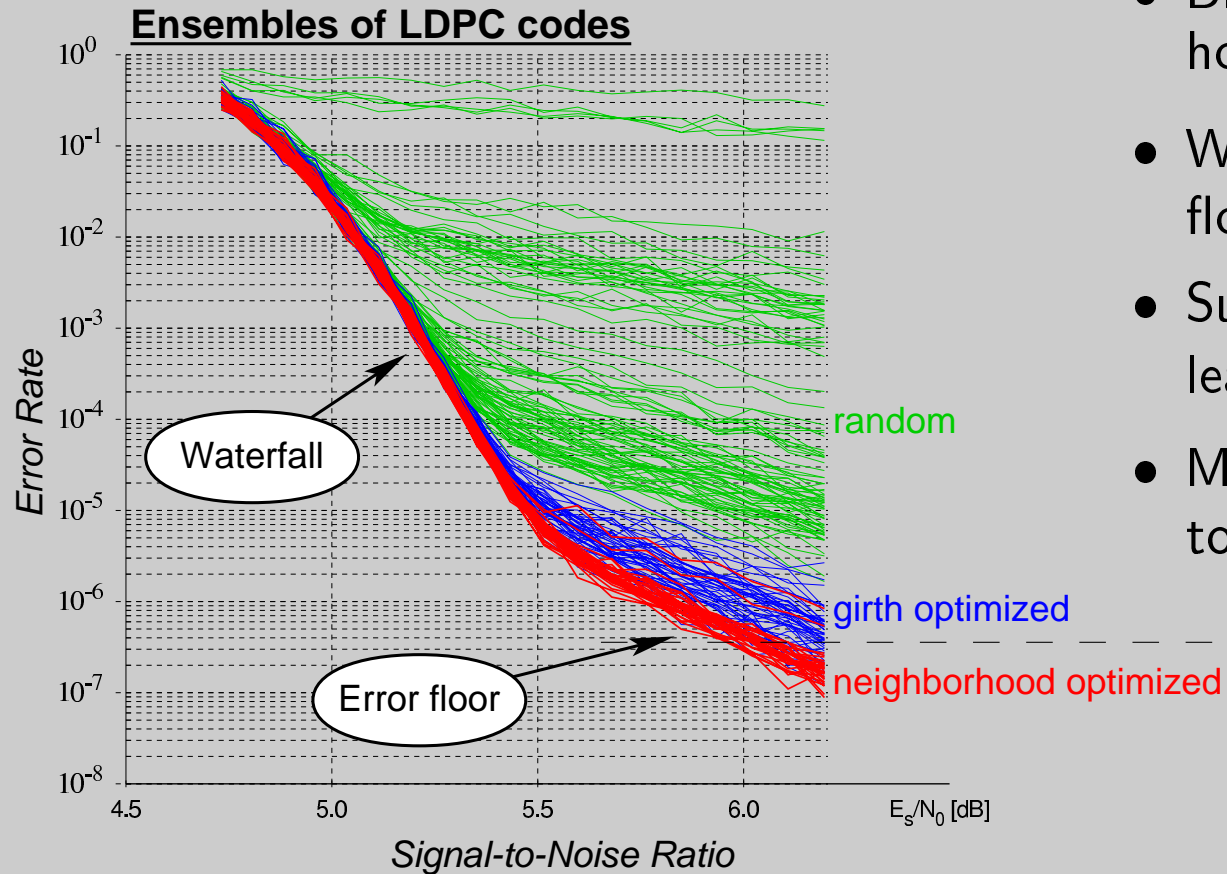
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Amoeba (downhill simplex method)



Numerical Recipes,
ch. 10, part 4

Error floors of LDPC codes



- BER *vs* SNR — how good is the code?
- Waterfall \leftrightarrow error floor transition
- Suboptimal decoding leads to error floor
- Monte Carlo fails to reach lowest BER

no go zone
too few errors for
Monte Carlo

Tom Richardson, Error floors of LDPC codes

Binary, linear error-correcting codes

$$\begin{pmatrix} 0 \\ 1 \\ 1 \\ 1 \\ 0 \\ \vdots \\ N_1 \\ \times \\ 1 \end{pmatrix} = \begin{pmatrix} \hat{G} & \begin{pmatrix} 1 & 0 & 1 \\ 1 & 1 & 0 \\ 0 & 1 & 1 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix} \\ N \times L \end{pmatrix} \begin{pmatrix} L \\ \times \\ 1 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \\ 1 \end{pmatrix}$$

codeword \vec{y} generator matrix information \vec{x}

$\vec{y} \rightarrow \vec{y} + \vec{\zeta}$
distortion of signal by channel

$$\begin{pmatrix} \hat{H} & \begin{pmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{pmatrix} \\ M \times N \end{pmatrix} \begin{pmatrix} N \\ \times \\ 1 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \\ 1 \\ 1 \\ 0 \\ \vdots \\ N_1 \end{pmatrix} = \begin{pmatrix} M \\ \times \\ 1 \end{pmatrix} \begin{pmatrix} 0 \\ 0 \\ 0 \\ \vdots \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \\ \vdots \\ 0 \end{pmatrix}$$

parity check matrix codeword \vec{y} syndrome vector

Decoding

$$\text{decoded codeword} = \underset{\text{all codewords}}{\operatorname{argmax}} \mathcal{P} \left(\text{codeword} \mid \begin{array}{l} \text{channel} \\ \text{output} \end{array} \right) \quad 2^{\#\text{bits}} \text{ operations}$$

Iterative decoding

(#edges) · (#iterations) operations

Checks vote for the bits value (unsatisfied check votes to flip the bit)

Proceed voting iteratively until convergence

$$\eta_{i\alpha}^{(n+1)} = h_i + \sum_{\beta \ni i} \sum_{\beta \neq \alpha} \tanh^{-1} \prod_{\substack{j \in \beta \\ j \neq i}} \tanh \eta_{j\beta}^{(n)}, \quad h_i = \frac{1}{2} \log \frac{p(y_i | +1)}{p(y_i | -1)}, \quad \eta_{i\alpha}^{(0)} \equiv 0$$

$$\eta_{i\alpha}^{(n+1)} = h_i + \sum_{\beta \ni i} \sum_{\beta \neq \alpha} \left(\prod_{\substack{j \in \beta \\ j \neq i}} \operatorname{sign} \eta_{j\beta}^{(n)} \right) \min_{\substack{j \in \beta \\ j \neq i}} \left| \eta_{j\beta}^{(n)} \right|$$

Decoding

$$\text{decoded codeword} = \underset{\text{all codewords}}{\operatorname{argmax}} \mathcal{P} \left(\text{codeword} \mid \text{channel output} \right) \quad 2^{\#\text{bits}} \text{ operations}$$

Iterative decoding

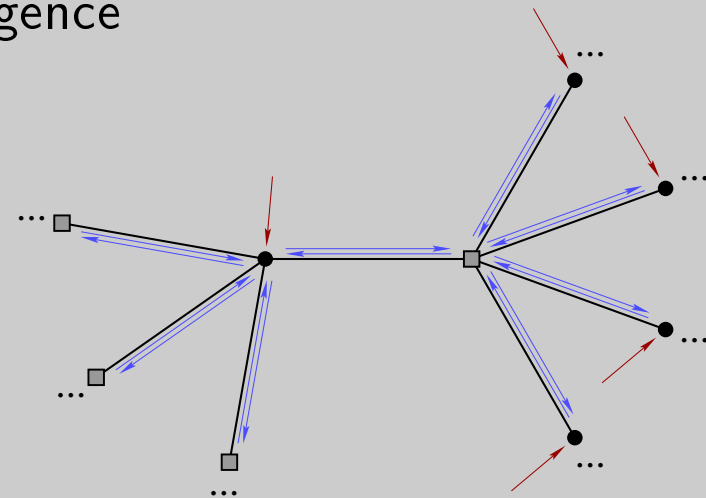
(#edges) · (#iterations) operations

Checks vote for the bits value (unsatisfied check votes to flip the bit)

Proceed voting iteratively until convergence

Message passing
Belief propagation on a graph

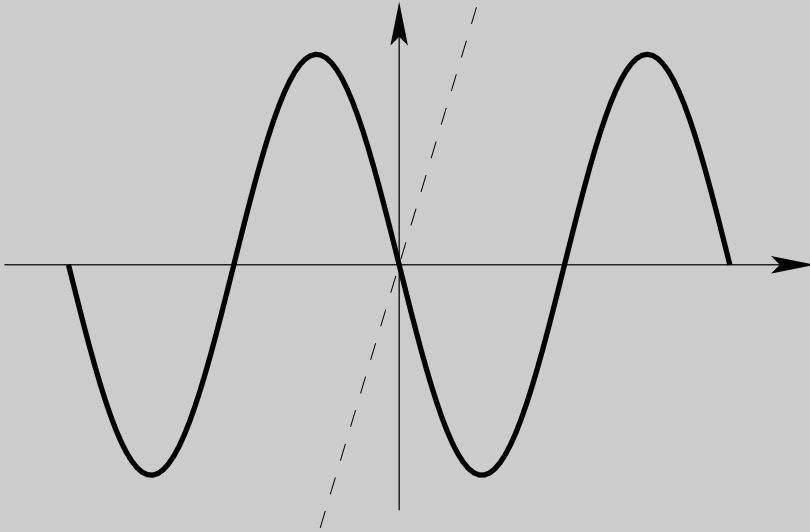
Linear programming



Iterative solution, works for trees

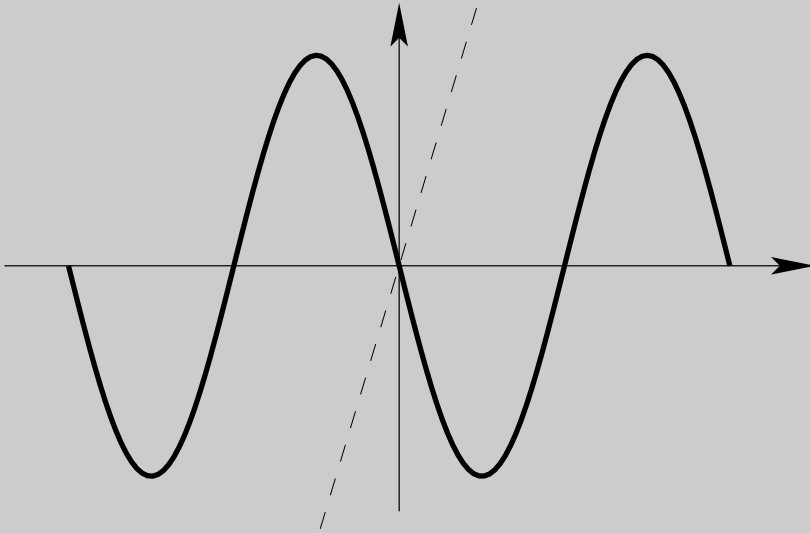
Bethe (1935), Peirls (1936), Gallager (1962), Pearl (1986), MacKay (1995)

Unstable iterations, period doubling

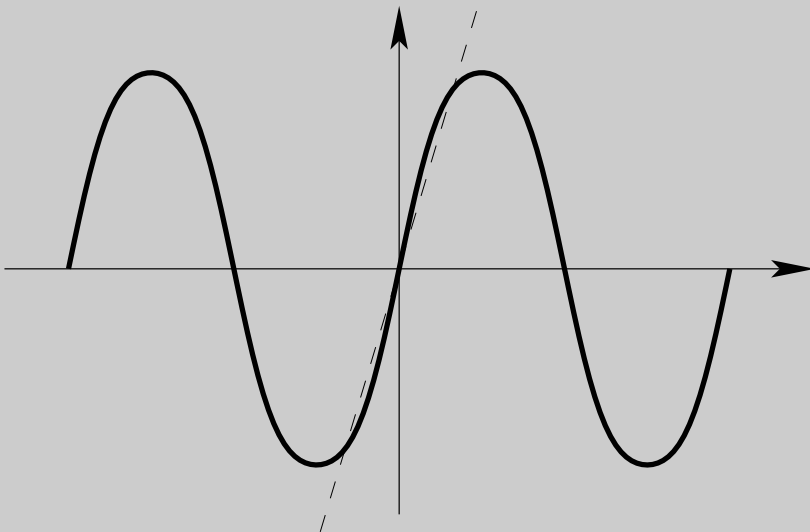


$$f(x) = -1.2 \sin(x)$$

Unstable iterations, period doubling

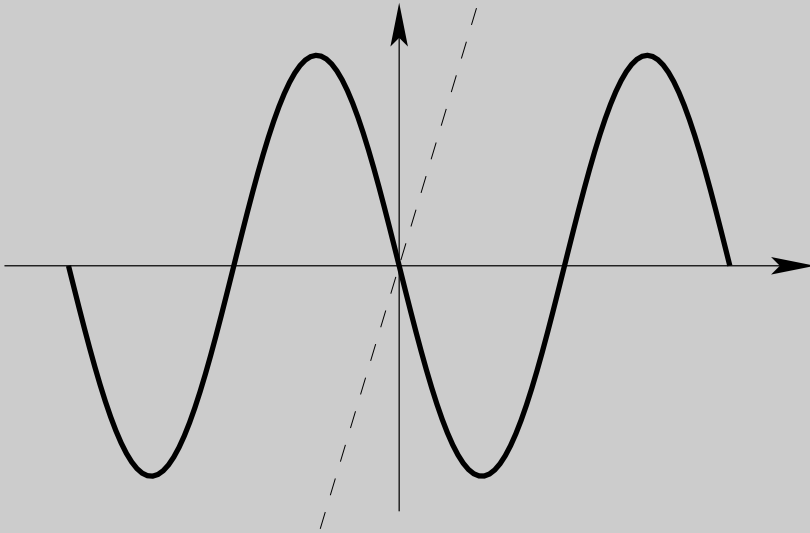


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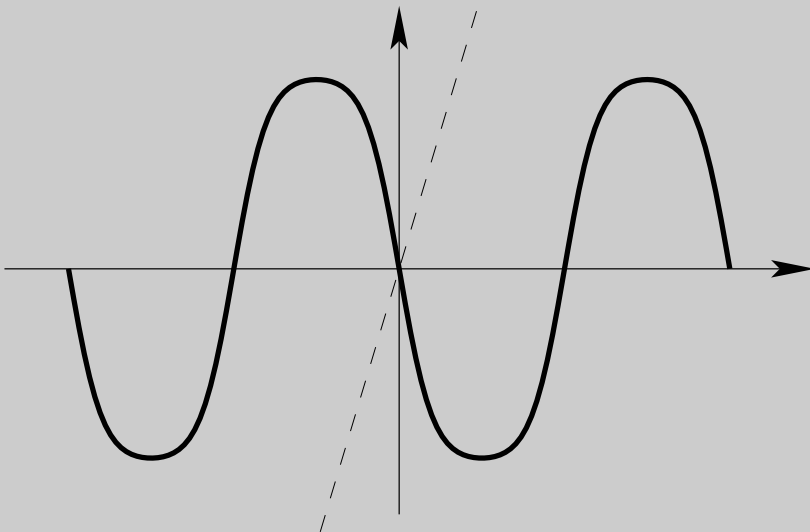


$$f(f(x))$$

Unstable iterations, period doubling

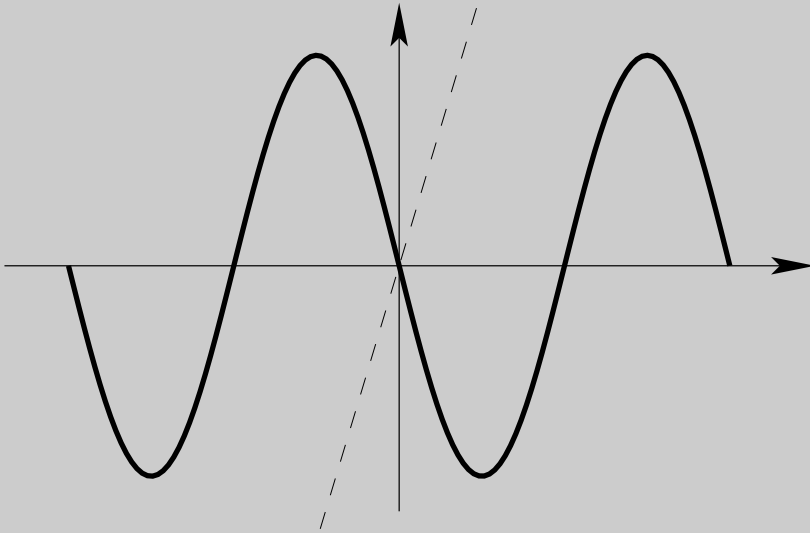


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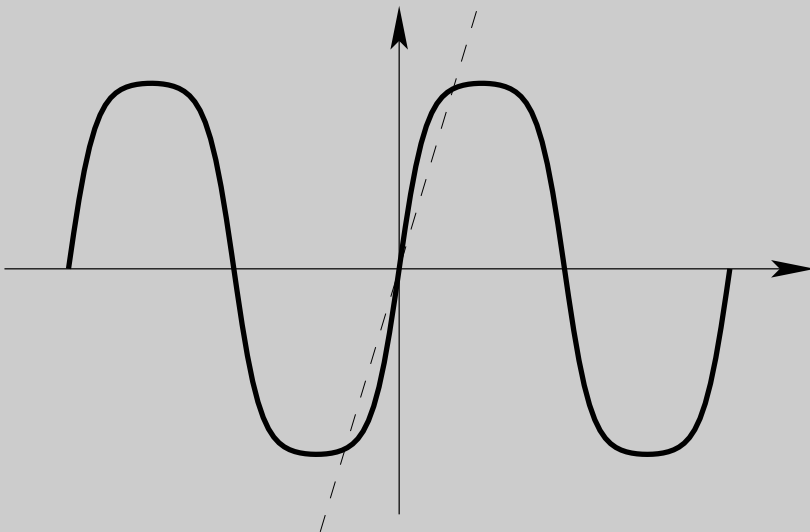


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Unstable iterations, period doubling

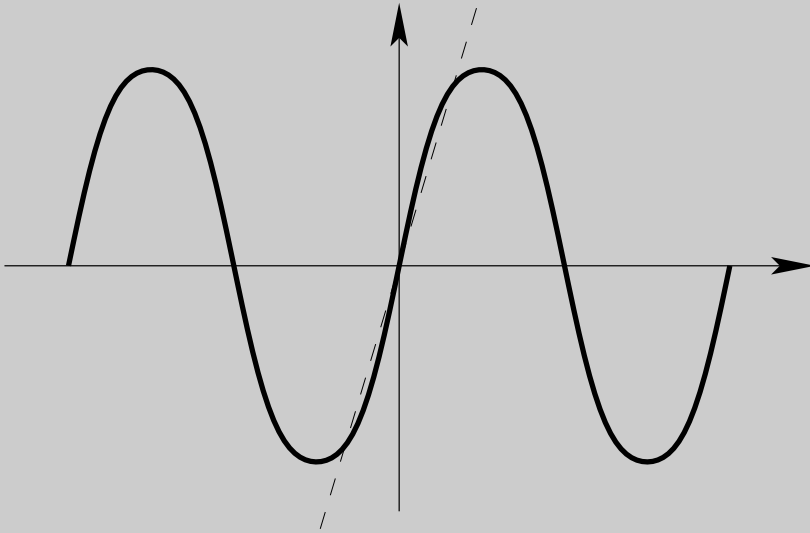


$$f(x) = -1.2 \sin(x)$$

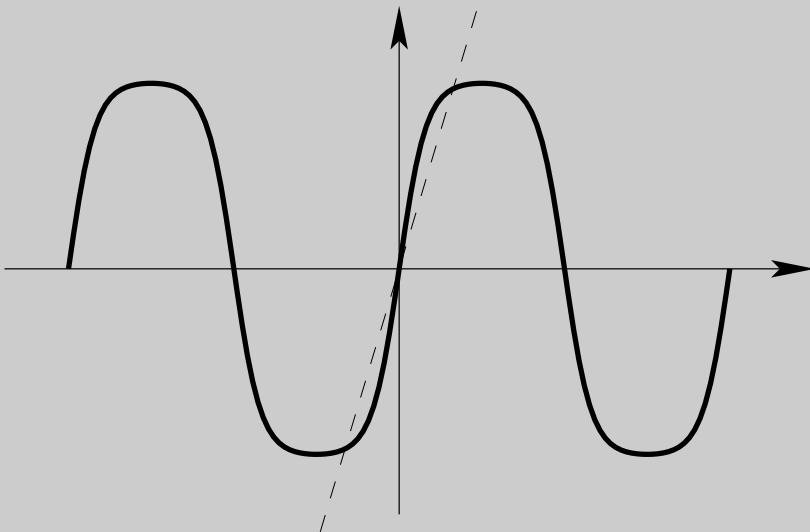


$$f(f(f(f(x))))$$

Unstable iterations, period doubling

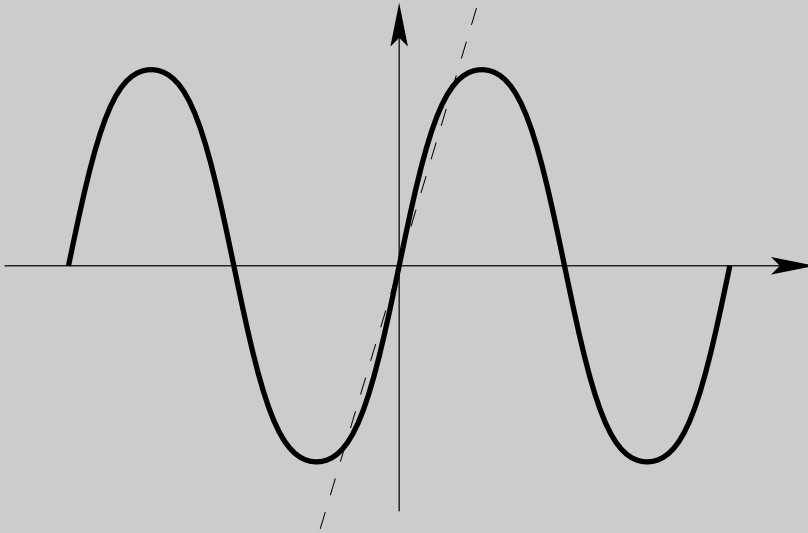


$$f(f(x))$$

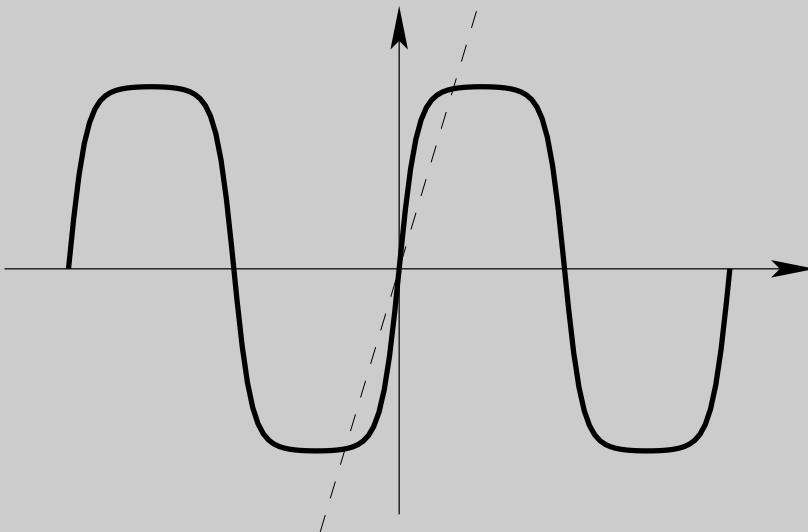


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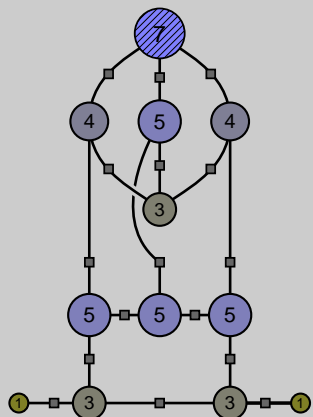


$$f(f(x))$$

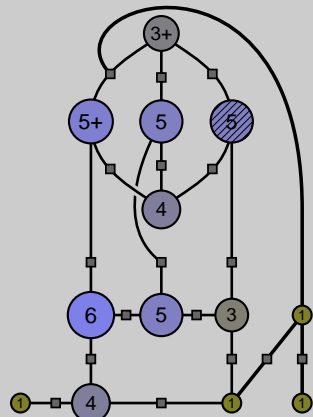


$$f(f(f(f(f(x))))))$$

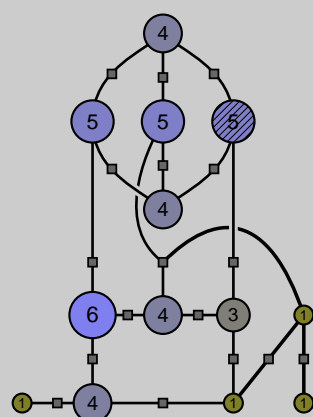
Instantons, Laplacian channel



$$l_{\text{ef}}^2 \approx 10.076$$

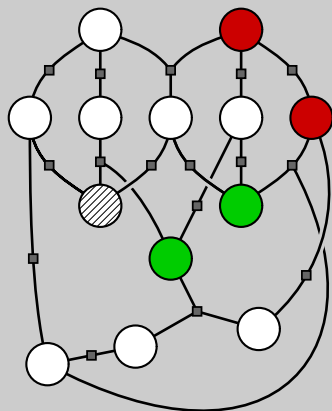


$$l_{\text{ef}}^2 \approx 10.203$$

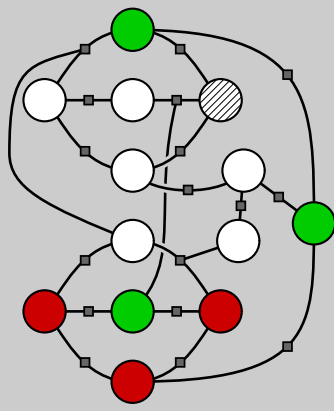


$$l_{\text{ef}}^2 \approx 10.298$$

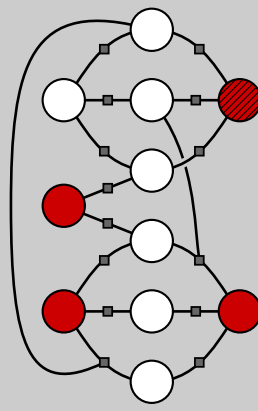
Gaussian



$$l_{\text{ef}} = 7.6$$



$$l_{\text{ef}} = 8.$$

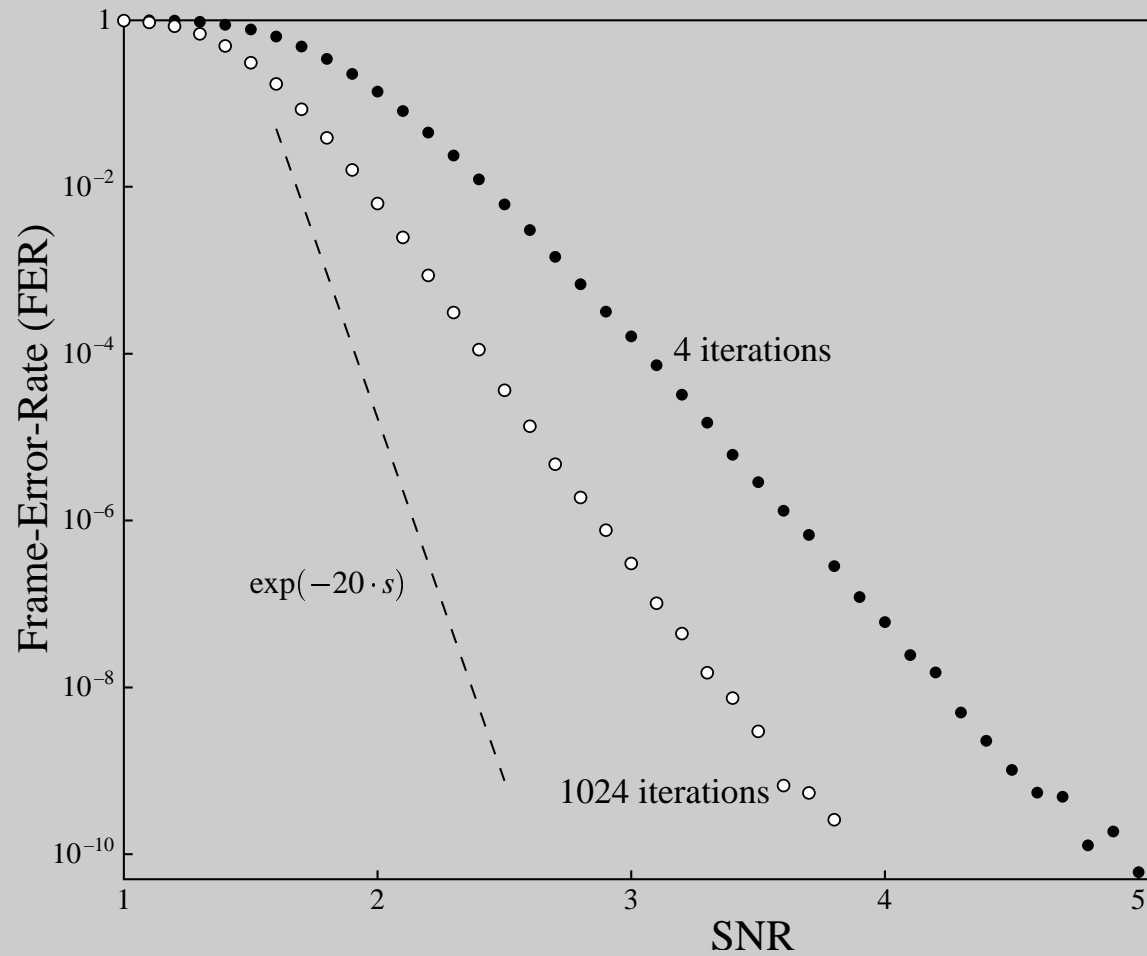


$$l_{\text{ef}} = 8.$$

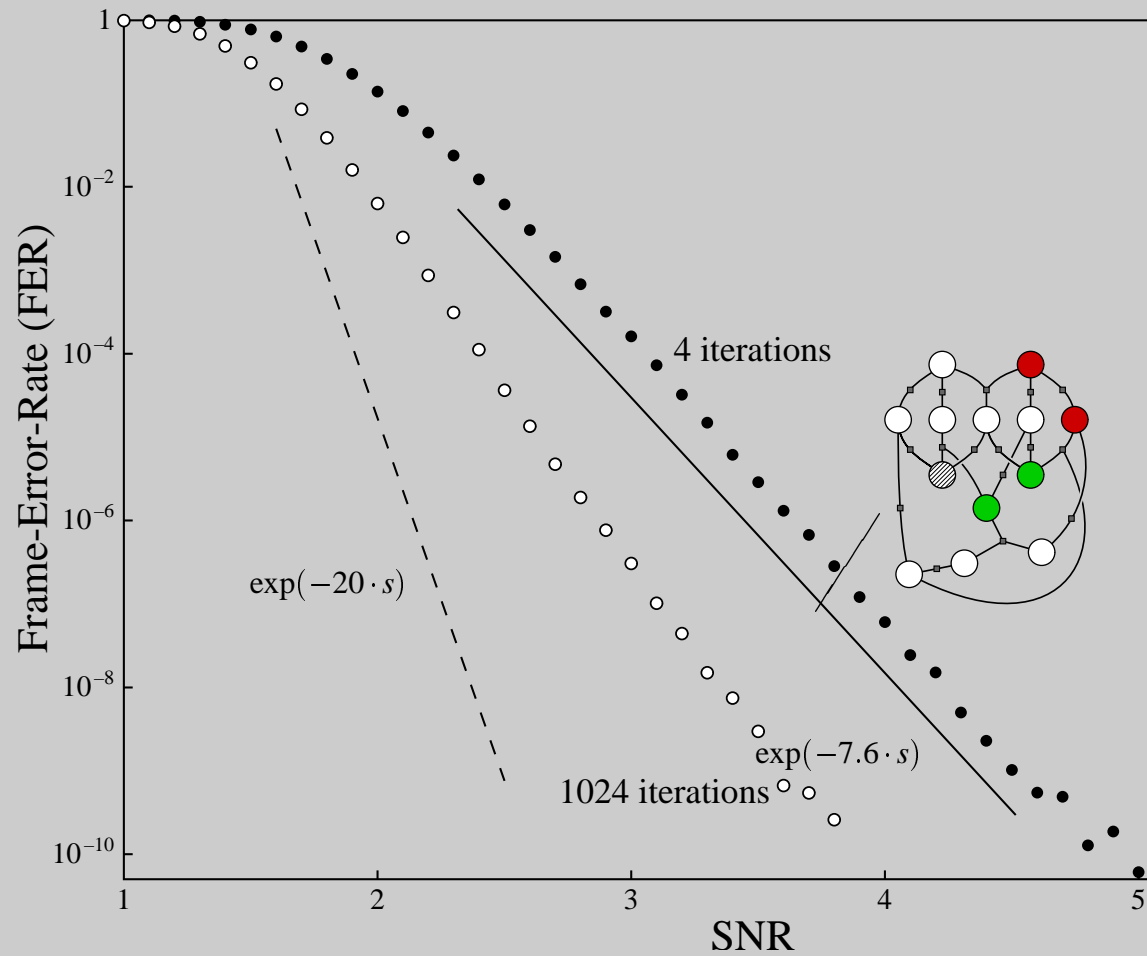
Laplacian
exponential tails

Stepanov, Chertkov,
2005 Allerton Conference [arxiv.org: cs.IT/0507031]

Frame-Error-Rate, Laplacian channel



Frame-Error-Rate, Laplacian channel



BP as minimization of Bethe free energy

Yedidia, Freeman, Weiss

Variables: beliefs $0 \leq b_i(\sigma_i), b_\alpha(\sigma_\alpha) \leq 1$

Conditions: normalization $\sum_{\sigma_i} b_i(\sigma_i) = \sum_{\sigma_\alpha} b_\alpha(\sigma_\alpha) = 1$

consistency $\sum_{\sigma_\alpha \setminus \sigma_i} b_\alpha(\sigma_\alpha) = b_i(\sigma_i)$

Function: Bethe free energy

$$\mathcal{F}_{\text{Bethe}} = - \sum_{\alpha} \sum_{\sigma_\alpha} b_\alpha(\sigma_\alpha) \ln f_\alpha(\sigma_\alpha) + \sum_{\alpha} \sum_{\sigma_\alpha} b_\alpha(\sigma_\alpha) \ln b_\alpha(\sigma_\alpha) - \sum_i (q_i - 1) \sum_{\sigma_i} b_i(\sigma_i) \ln b_i(\sigma_i)$$

$$f_\alpha(\sigma_\alpha) \equiv \exp \left(\sum_{i \in \alpha} h_i \sigma_i / q_i \right) \delta \left(\prod_{i \in \alpha} \sigma_i, 1 \right)$$

Minimization \implies BP equation: $\eta_{i\alpha} = h_i + \sum_{\beta \ni i, \beta \neq \alpha} \tanh^{-1} \left(\prod_{j \in \beta} \tanh \eta_{j\beta} \right)$

Relaxed (smoothed) decoding

$$\mathcal{L} = \mathcal{F}_{\text{Bethe}} + (\text{Lagrangian multipliers}) \cdot (\text{conditions})$$

Minimization: $\frac{\delta \mathcal{L}}{\delta(\text{beliefs})} = 0, \quad \frac{\delta \mathcal{L}}{\delta(\text{Lagrangian multipliers})} = 0$

Iterative scheme (BP): $\eta_{i\alpha}^{(n+1)} = h_i + \sum_{\beta \ni i}^{\beta \neq \alpha} \tanh^{-1} \left(\prod_{j \in \beta} \tanh \eta_{j\beta}^{(n)} \right)$

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tree-based re-parametrization

Wainwright, Jaakola, Willsky

concave-convex procedure

Yuille

Heskes, Albers, Kappen

Relaxed (smoothed) decoding

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Dynamics: $\frac{d(\text{beliefs})}{d\tau} = -\frac{\delta \mathcal{L}}{\delta(\text{beliefs})}$
 $\frac{d(\text{Lagrangian multipliers})}{d\tau} = -\frac{\delta \mathcal{L}}{\delta(\text{Lagrangian multipliers})}$

Relaxed (smoothed) decoding

$$\mathcal{L} = \mathcal{F}_{\text{Bethe}} + (\text{Lagrangian multipliers}) \cdot (\text{conditions})$$

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$\frac{d(\text{Lagrangian multipliers})}{d\tau} = -\frac{\delta \mathcal{L}}{\delta(\text{Lagrangian multipliers})}$

$\eta_{i\alpha}^{(n+1)} + \frac{1}{\Delta} \sum_{\beta \ni i} \eta_{i\beta}^{(n+1)} = h_i + \sum_{\beta \ni i}^{\beta \neq \alpha} \tanh^{-1} \left(\prod_{j \in \beta} \tanh \eta_{j\beta}^{(n)} \right) + \frac{1}{\Delta} \sum_{\beta \ni i} \eta_{i\beta}^{(n)}$

Relaxed (smoothed) decoding

Iterative scheme (BP): $\eta_{i\alpha}^{(n+1)} = h_i + \sum_{\substack{\beta \neq \alpha \\ \beta \ni i}} \tanh^{-1} \left(\prod_{j \in \beta} \tanh \eta_{j\beta}^{(n)} \right)$

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Smoothed (relaxed, damped) decoding

Iterative scheme (BP): $\eta_{i\alpha}^{(n+1)} = h_i + \sum_{\beta \ni i, \beta \neq \alpha} \tanh^{-1} \left(\prod_{j \in \beta} \tanh \eta_{j\beta}^{(n)} \right)$

$$\eta_{i\alpha}^{(n+1)} + \frac{1}{\Delta} \sum_{\beta \ni i} \eta_{i\beta}^{(n+1)} = h_i + \sum_{\beta \ni i, \beta \neq \alpha} \tanh^{-1} \left(\prod_{j \in \beta} \tanh \eta_{j\beta}^{(n)} \right) + \frac{1}{\Delta} \sum_{\beta \ni i} \eta_{i\beta}^{(n)}$$