# Leadership Statistics in Random Structures 

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#### Abstract

The largest component ("the leader") in evolving random structures often exhibits universal statistical properties. This phenomenon is demonstrated analytically for two ubiquitous structures: random trees and random graphs. In both cases, lead changes are rare as the average number of lead changes increases quadratically with logarithm of the system size. As a function of time, the number of lead changes is self-similar. Additionally, the probability that no lead change ever occurs decays exponentially with the average number of lead changes.


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Extreme statistics are important in science, engineering, and society as they dictate catastrophic events, systems robustness, financial indices, etc. The theory of extreme statistics provides a powerful analysis framework and prediction tool [1, 2]. However, it is limited to ensembles of independent random variables. Even though most practical applications involve correlated random variables, such cases remain largely unexplored [3-5].

We investigate extremal characteristics of two basic random structures: random trees and random graphs. Random trees appear in data storage algorithms in computer science [6-8] and in physical processes such as collisions in gases [9]. Random graphs [10, 11] have numerous applications to theoretical computer science $[7,8]$, social networks $[12,13]$, and physical processes such as polymerization [14].

We focus on the largest component, the leader, and ask: What is the size of the leader? How does the number of lead changes depend on time and system size? What is the probability that the leader never changes? Similar questions were investigated in growing networks [15, 16], and related leadership statistics were studied in random graphs by Erdős and Łuczak [17].

Random trees and random graphs are special cases of aggregation processes and hence, we analyze them using the rate equation approach [18-21]. Characterization of leadership statistics is sensible only for finite systems. We thus consider large yet finite systems for which the rate equation approach yields the leading asymptotic dependence on the system size [22-25].

Our main result is that the total number of lead changes $L$ grows as $L(N) \sim[\ln N]^{2}$ with the system size $N$. This as well as other leadership statistics are universal as they characterize both random trees and random graphs. The time dependent number of lead changes $L(t, N)$ attains the scaling form $(\ln N)^{2} F(x)$ with the

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FIG. 1: Illustration of the tree merger process with system size $N=4$. Nodes are denoted by circles, with the top node (the root) denoted by a bullet. The bottom nodes are the leafs.
scaling variable $x=\ln k_{*} / \ln N$ where $k_{*}$ is the typical component size at time $t$. Additionally, the probability that no lead change has ever occurred decays as $e^{-L}$.

We start with the simpler case of random trees. A tree is a hierarchical structure of connected nodes. We term the top node the root and the bottom nodes the leafs (see Fig. 1). Initially, the system consists of $N$ trivial trees in which the root is also a leaf. We grow an ensemble (a forest) of trees using a recursive procedure in which two trees are picked at random and attached to a common root (Fig. 1). This merging process is repeated until a single tree with $N$ leafs emerges.

We treat the merger process dynamically. Let $n$ be the number of trees. The transition $n \rightarrow n-1$ occurs with rate $r_{n}$ proportional to the total number of pairs. Choosing $2 / N$ as the merger rate for each pair (i.e., $\left.r_{n}=n(n-1) / N\right)$ is convenient as in the thermodynamic limit $N \rightarrow \infty$, the normalized density $c=n / N$ evolves according to $\frac{d}{d t} c=-c^{2}$. Subject to the initial condition $c(0)=1$ the density is $c(t)=(1+t)^{-1}$. We state the results in terms of time $t$; the results can be easily reformulated in terms of the number of trees $n=N(1+t)^{-1}$.

The size distribution is obtained similarly. Let $n_{k}$ be the number of trees with $k$ leafs. The normalized density $c_{k}=n_{k} / N$ evolves according to the Smoluchowsky rate equation $\frac{d}{d t} c_{k}=\sum_{i+j=k} c_{i} c_{j}-2 c c_{k}$ with the monodisperse initial conditions $c_{k}(0)=\delta_{k, 1}$. The rate equation reflects the fact that trees are merged randomly, independent of their size. It can be solved (using the generating
functions technique for example) to give $[18,19]$

$$
\begin{equation*}
c_{k}(t)=\frac{t^{k-1}}{(1+t)^{k+1}} \tag{1}
\end{equation*}
$$

In the long time limit, the size distribution attains a simple self-similar behavior

$$
\begin{equation*}
c_{k}(t) \simeq k_{*}^{-2} \Phi\left(k / k_{*}\right), \quad \Phi(x)=e^{-x} \tag{2}
\end{equation*}
$$

with the typical size $k_{*} \simeq t$.
What is the average size of the largest tree (the leader)? Using the size distribution, we can answer an even more general question. Let $l_{r}(t)$ be the average size of the $r$ largest tree with the leader $l \equiv l_{1}$. From the cumulative distribution $u_{k}=\sum_{j \geq k} n_{j}=N t^{-1}[t /(1+t)]^{k}$ and the relation $u_{l_{r}}=r$, the size of the $r$ th leader is

$$
\begin{equation*}
l_{r}(t, N)=\frac{\ln [N / r t]}{\ln [(1+t) / t]} \tag{3}
\end{equation*}
$$

There are two regimes of behavior. In the short time limit, $t \ll 1$, one has $l_{r}(t, N)=1+\ln [N / r] / \ln [1 / t]$. Moreover, from $n_{k} \simeq N t^{k-1}$ the first dimer appears at $t_{2} \simeq N^{-1} ;$ the first trimer appears at $t_{3} \simeq N^{-1 / 2}$ and then, there are of the order $N^{1 / 2}$ dimers, so this trimer results from the leading dimer with probability of the order $N^{-1 / 2}$. At the crossover point, $t \approx 1$, the size of the leader varies logarithmically with the system size, $l(t \approx 1, N) \sim \ln N$. In the long time limit, $t \gg 1$, the size of the leader grows linearly (up to a logarithmic correction) with time

$$
\begin{equation*}
l_{r}(t, N) \simeq t \ln \frac{N}{r t} \tag{4}
\end{equation*}
$$

What is the average size $l_{w}$ of the winner (the leader that never relinquishes the lead)? At what time $t_{w}$ does the winner emerge? Both quantities grow linearly with $N: l_{w} \simeq \alpha N$ and $t_{w} \simeq \beta N$ as follows from Eq. (4). The curve $\alpha=-\beta \ln \beta$ has an extremum at $\alpha=\beta=$ $e^{-1} \cong 0.36788$ thereby implying the bounds: $\alpha, \beta<e^{-1}$. Finding the values of $\alpha$ and $\beta$ is an open problem.

How many lead changes $L(t, N)$ occur as a function of time? As a function of system size? What is the total number of lead changes $L(N) \equiv L(t=\infty, N)$ ? In our definition, a lead change occurs when two trees (none of which is the leader) merge to become larger than the leader. For short times, $t \ll 1$, almost every lead change introduces a new leader, and since the leader grows in increments of unity $L(t, N)=l(t, N)-1$. For long times, $t \gg 1$, consider the cumulative distribution $u_{k} \simeq N t^{-1} \exp (-k / t)$. Its growth rate immediately gives the rate by which the leader is surpassed, $\frac{d}{d t} L(t, N)=$ $\left.\frac{d}{d t} u_{k}\right|_{k=l}$ (this rate properly accounts for events in which two non-leaders are merging to surpass the leader). As $u_{l} \simeq 1$, we have $\frac{d}{d t} L(t, N) \simeq l t^{-2} \simeq t^{-1} \ln \frac{N}{t}$ from which the time-dependent number of lead changes is

$$
\begin{equation*}
L(t, N) \simeq \ln t \ln N-\frac{1}{2}(\ln t)^{2} \tag{5}
\end{equation*}
$$



FIG. 2: The normalized time dependence of the number of lead changes for random trees, $L(t, N) / L(N)$, versus the scaling variable $x=\ln t / \ln N$. The simulation data, representing an average over $10^{3}$ Monte Carlo runs, is compared with the theoretical prediction $2 F(x)=2 x-x^{2}$.

Interestingly, this quantity obeys the scaling form

$$
\begin{equation*}
L(t, N)=(\ln N)^{2} F(x), \quad x=\frac{\ln t}{\ln N} \tag{6}
\end{equation*}
$$

with the quadratic scaling function: $F(x)=x-\frac{1}{2} x^{2}$ (Fig. 2). The scaling variable is unusual: a ratio of logarithms, in contrast with the ordinary ratio underlying the size distribution (2).

To check these theoretical predictions, we performed large-scale Monte Carlo simulations. In the simulations, randomly chosen trees are merged repeatedly. Keeping track of the leader and averaging over many independent realizations, we observe a scaling behavior that is consistent with Eq. (6). However, as a function of the system size, the convergence is slow because the scaling variable involves logarithms.

The time dependent behavior can be used to obtain the total number of lead changes. Substituting $t_{w}=\beta N$ into Eq. (5) gives

$$
\begin{equation*}
L(N) \simeq A(\ln N)^{2} \tag{7}
\end{equation*}
$$

with $A=F(1)=1 / 2$. Both the leading asymptotic behavior and the $\ln N$ correction are confirmed numerically. Moreover, the numerical prefactor $A=0.50(1)$ agrees with the theoretical prediction (Fig. 3). Since $L(t \approx 1, N) \sim \ln N$, the majority of lead changes occur when $t \gg 1$.

How is the number of lead changes distributed? What is the probability that no lead change occur? Let $P_{n}(t, N)$ be the probability that $n$ lead changes occur till time $t$. We assume that the flux surpassing the leader characterizes the evolution of the probability distribution $\frac{d}{d t} P_{n}=\left(\frac{d}{d t} L\right)\left[P_{n-1}-P_{n}\right]$. With the initial condition


FIG. 3: The total number of lead changes $L(N)$ versus the system size $N$. Shown are simulation results for Random Trees (RT) and Random Graphs (RG) representing an average over $10^{4}$ realizations.
$P_{n}(0, N)=\delta_{n, 0}$, the distribution is Poissonian

$$
\begin{equation*}
P_{n}(t, N)=\frac{[L(t, N)]^{n}}{n!} e^{-L(t, N)} \tag{8}
\end{equation*}
$$

Consequently, the distribution of the total number of lead changes $P_{n}(N) \equiv P_{n}(t=\infty, N)$ is also Poissonian: $P_{n}(N)=\frac{L^{n}}{n!} e^{-L}$ with $L \equiv L(N)$ given by (7). Hence, the standard deviation in the number of lead changes $\sigma(N)$ grows as $\sigma(N) \simeq \sqrt{A} \ln N$. Furthermore, the probability that no lead change occur (the survival probability of the first leader) $S(N) \equiv P_{0}(N)$ decays faster than a power-law but slower than a stretched exponential

$$
\begin{equation*}
S(N)=\exp [-L] \simeq \exp \left[-A(\ln N)^{2}\right] \tag{9}
\end{equation*}
$$

The asymptotic $N$-dependence is confirmed numerically (Fig. 4).

We now consider graphs, grown randomly as follows. Initially, the system consists of $N$ single-node graphs. Then, two nodes are picked at random and a link is drawn between them. If they belong to two distinct graphs, the two become one. This process is repeated indefinitely. Let $n_{k}$ be the number of graphs of size $k$. The normalized density $c_{k}=n_{k} / N$ evolves according to the rate equation $\frac{d}{d t} c_{k}=\frac{1}{2} \sum_{i+j=k} i j c_{i} c_{j}-k c_{k}$ with the monodisperse initial conditions $c_{k}(0)=\delta_{k, 1}$. The rate equation reflects the fact that two graphs are connected with rate proportional to the product of their sizes. This equation can be solved using generating functions to give [20, 21]

$$
\begin{equation*}
c_{k}(t)=\frac{(k t)^{k-1}}{k \cdot k!} e^{-k t} \tag{10}
\end{equation*}
$$

At time $t=1$, the system undergoes a gelation transition: it develops a giant component that eventually takes


FIG. 4: The survival probability of the initial leader $S(N)$ versus the system size $N$. The number of realizations was $10^{10}$ and $10^{8}$ for random trees and random graphs, respectively.
over the entire system. Close to the gelation time, the size distributions attains the scaling behavior

$$
\begin{equation*}
c_{k}(t) \simeq k_{*}^{-5 / 2} \Phi\left(k / k_{*}\right), \quad \Phi(x) \propto x^{-5 / 2} e^{-x / 2} \tag{11}
\end{equation*}
$$

The typical size diverges: $k_{*} \simeq(1-t)^{-2}$ as $t \rightarrow 1$. At the gelation time, the size distribution develops an algebraic tail $c_{k}(t=1) \sim k^{-5 / 2}$. Hence, the cumulative distribution is $u_{k} \sim N k^{-3 / 2}$ and with the criterion $u_{l_{w}} \sim 1$ we obtain the average size of the giant component (the winner): $l_{w} \sim N^{2 / 3}$ [10]. The time at which it emerges is $1-t_{w} \sim N^{-1 / 3}$.

The size of the leader is again estimated from the relation $u_{l}=1$. For $t \ll 1$, the size of the leader $l(t, N)$, the number of lead changes $L(t, N)$, as well as the number of distinct leaders are all approximately equal and the same as for random trees. The number of lead changes is of the order $\ln N$ in this phase; furthermore, $L(t, N) \sim \ln N$ for $t<1$. The behavior near the gelation time $1-t \ll 1$ is a bit more interesting. From the large- $k$ behavior, $u_{k} \sim N(1-t)^{-2} k^{-5 / 2} \exp \left[-k(1-t)^{2} / 2\right]$, and $u_{l}=1$ we arrive at following implicit relation for the leader $l(t, N)$ : $l \simeq 2(1-t)^{-2} \ln N-3(1-t)^{-2} \ln l$. Inserting the zeroth order approximation $l^{(0)}=2(1-t)^{-2} \ln N$ into $\ln l$ on the right-hand side of the above relation and ignoring $\ln \ln N$ terms yields the leader size

$$
\begin{equation*}
l \simeq \frac{2}{(1-t)^{2}} \ln \left[N(1-t)^{3}\right] . \tag{12}
\end{equation*}
$$

The rate by which the leader changes is estimated from $\frac{d}{d t} L=\left.\frac{d}{d t} u_{k}\right|_{k=l} \simeq l(1-t)$. Substituting Eq. (12) and integrating, the number of lead changes is

$$
\begin{equation*}
L(t, N) \simeq 2 \ln N \ln \frac{1}{1-t}-3\left[\ln \frac{1}{1-t}\right]^{2} \tag{13}
\end{equation*}
$$

It attains the scaling form

$$
\begin{equation*}
L(t, N) \simeq(\ln N)^{2} F(x), \quad x=\frac{\ln \frac{1}{1-t}}{\ln N} \tag{14}
\end{equation*}
$$

with the scaling function $F(x)=2 x-3 x^{2}$. As this behavior holds up to time $t_{w}$, where $1-t_{w} \sim N^{-1 / 3}$, the total number of lead changes grows according to Eq. (7) with $A=F(1 / 3)=1 / 3$. Furthermore, the distribution of lead changes is Poissonian as in (8) and the survival probability decays according to (9).

As terms of the order $\ln \ln N / \ln N$ were neglected, the scaling behavior and the leading asymptotic behavior may be realized only for extremely large systems. Moreover, the computational cost of random graph simulations is larger because graphs are chosen with probability proportional to their size. Nevertheless, we can confirm the predicted system size dependence of $L(N)$ (Fig. 2) and $S(N)$ (Fig. 3) numerically. The prefactor $A=0.20(2)$ is lower than the theoretical value $A=1 / 3$, perhaps due to the slow convergence.

Let us compare random trees and random graphs. They seem very different, e.g., the gelation transition occurs in one case but not in the other. Yet, they exhibit similar extremal characteristics. In both cases, the total number of lead changes $L(N)$ grows as $[\ln N]^{2}$ and the survival probability decays as $\exp [-L]$. Moreover, even the seemingly distinct temporal characteristics can be reconciled, e.g., in both cases the size distribution attains the scaling form $c_{k}(t) \propto \Phi\left(k / k_{*}\right)$ when $t \rightarrow \infty$ (for random trees) and $t \rightarrow 1$ (for random graphs). Of course, the actual time dependence of the typical scale is different: $k_{*} \sim t$ and $k_{*} \sim(1-t)^{-2}$, respectively. The size of the leader can be rewritten as $l \approx k_{*} \ln \left[N / k_{*}^{\gamma}\right]$ with $\gamma=1$ and $3 / 2$, respectively. Furthermore, Eqs. (6) and (14) can be reconciled by writing the scaling variable in the unified form $x=\ln k_{*} / \ln N$.

We now restrict our attention to the survival probability $S(N)$ and supplement the leading behavior (9) with bounds. Consider random trees. A lower bound for $S(N)$ is obtained from a greedy scenario in which all merger events involve the leader till it reaches size $N / 2$. The probability that the second merger involves the leading dimer is $(N-1)^{-1}$; the probability that the third merger involves the leading trimer is $(N-2)^{-1}$; etc. Thus, the greedy scenario is realized with probability $\prod_{j<N / 2} \frac{1}{N-j}$, thereby providing the lower bound $S(N)>\frac{(N / 2)!}{(N-1)!}$. An upper bound can be obtained by estimating the number of trees the size of the leader. There are of the order $N^{1 / 2}$ dimers when the first trimer is born, $n_{2}\left(t_{3}=N^{-1 / 2}\right)=N^{1 / 2}$. The leading dimer retains the lead with probability inversely proportional to the number of dimers, $N^{-1 / 2}$. Similarly, this leading trimer retains the lead with probability proportional to $N^{-1 / 3}$. Therefore the upper bound $\prod_{j<\ln N} N^{-1 / j}$ is estimated as $N^{-\ln \ln N}$ (the cutoff $j<\ln N$ is dictated by the size of the leader at the crossover time $t \approx 1$ ). Hence, the survival probability obeys

$$
\begin{equation*}
\left(\frac{e}{2 N}\right)^{\frac{N}{2}}<S(N)<\exp [-(\ln N) \cdot(\ln \ln N)] \tag{15}
\end{equation*}
$$

where the lower bound was simplified using the Stirling formula. Note that the upper bound merely assures that
the lead never changes in the early phase $t<1$ when the average number of lead changes is only $\ln N$.

For random graphs, the greedy scenario is again simple to analyze since the probability that in a system with a leader of size $j$ and $N-j$ monomers the probability that the next merger involves the leader is $p_{j}=[j(N-j)] /\left[j(N-j)+\frac{1}{2}(N-j)(N-j-1)\right]$ or $p_{j}=(2 j) /(N+j-1)$. The product $\prod_{j<N / 2} p_{j}$ provides the lower bound. Asymptotically, the lower bound decays as $\lambda^{N}$ with $\lambda=(2 / 3)^{3 / 2}=0.544331 \ldots$. On the other hand, repeating the above argument yields the same upper bound, so

$$
\begin{equation*}
\lambda^{N}<S(N)<\exp [-(\ln N) \cdot(\ln \ln N)] \tag{16}
\end{equation*}
$$

The upper bound is again much closer to the actual asymptotic behavior.

In conclusion, random graphs and random trees exhibit similar leadership characteristics. As in random growing networks [16], lead changes are infrequent given that the overall number of lead changes increases only logarithmically with the system size. The time dependent number of lead changes approaches a self-similar form asymptotically. The convergence to the asymptotic behavior is much slower for extremal statistics compared with size statistics due to the various logarithmic dependences. Consequently, the asymptotic behavior may difficult to detect in practice, especially for random graphs.

To obtain the extremal characteristics, we employed the scaling behavior of the size distribution outside the scaling regime, namely, at sizes much larger than the typical size where, at least formally, statistical fluctuations can no longer be ignored. Interestingly, the resulting system size dependence for the various leadership statistics appears to be asymptotically exact. Further analysis is needed to illuminate the role of statistical fluctuations, for example by characterizing corrections to the leading behavior.

The virtue of the rate equation approach to analyzing extremal characteristics is its simplicity, robustness, and generality. It applies to general aggregation processes where the merger rate may depend in a complicated manner on the aggregate size or in situations where there is an underlying spatial structure. We find that the above leadership statistics extend to algebraic merger rates as well as to aggregation in one spatial dimension [26]. This method is also applicable to other extremal features including for example laggard (smallest component) statistics. In the case of random trees, for instance, the total number of laggard changes grows logarithmically with the system size.

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