Operations-Based Planning for Placement and Sizing of Energy Storage in a Grid With a High Penetration of Renewables

K. Dvijotham¹ M. Chertkov² S. Backhaus³

¹Computer Science and Engineering University of Washington, Seattle

²Center for Nonlinear Studies and Theoretical Division Los Alamos National Laboratory

³Material Physics and Applications Division Los Alamos National Laboratory

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The US Power Grid : The Present



"... the greatest engineering achievement of the 20th century."

[National Academy of Engineering '10]

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The Grid : Operation and Planning

- Expansion/Upgrade Planning : Adding new resources (lines/generators etc.) to the grid
- Operation : How to dispatch generation given current state of grid and forecasted loads/renweables
- Presently Decoupled

The grid is changing:

- Iarge number of distributed power sources
- increasing adoption of renewables
- ⇒ large-scale, complex, & heterogeneous networks with stochastic disturbances



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Implications

- Fluctuations in generation now become significant
- Transmission lines overload, Generators hit limits etc.
- Need "smart control" to compensate
- Planning and Operation can no longer be decoupled

Our Solution

- Use Energy Storage to mitigate renewable fluctuations
- Use Optimal Control to decide how to dispatch power from storage
- Use resulting optimal solutions on a large number of scenarios (profiles of renewable fluctuations) to decide optimal sizing and placement of energy storage

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Sneak Peak at Results

- With small amounts of Energy Storage, can handle 30% Renewable Penetration
- Energy Storage only required at small number of nodes
- Using Operations information in Planning/Sizing of Energy Storage yields best results

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Power Grid : Description

- Topology specified by graph $\mathcal{G} = (\mathcal{V}, \mathcal{E}), |\mathcal{V}| = n$.
- Nodes *i* denote generators/loads, Edges (*i*, *j*) denote transmission lines
- Neb(*i*) denotes the neighbors of node *i* in the grid

Three kinds of nodes in the grid:

- Controllable Generators G_g: Output Controllable within Limits
- Renewable Generators *G_r*:
 Output Fluctuating over time
- Loads G_I : Assumed Fixed Here

Modified RTS-96 Grid





Given power injected/consumed at every node, compute how much power flows through each transmission line.

- Power generated/consumed at node *i* : **p**_{*i*}
- Power flowing from node *i* to *j*: $\mathbf{f}_{(i,j)}$
- Complex power phase at node $i: \psi_i$
- Inductance on line (i, j): $\mathbf{x}_{(i,j)}$

DC Power Flow Equations

$$\mathbf{p}_i = \sum_{j \in \text{Neb}(i)} \frac{\psi_j - \psi_i}{\mathbf{x}_{(i,j)}}$$
$$\mathbf{f}_{(i,j)} = \mathbf{x}_{(i,j)}(\psi_j - \psi_i),$$

Constraints on Power Generation

- Graph Laplacian $\mathbf{L} \in \mathbf{R}^{|\mathcal{V}| \times |\mathcal{V}|}$, $\mathbf{L}_{ii} = -\sum_{j \in \text{Neb}(i)} \frac{1}{\mathbf{x}_{(i,j)}}$, $\mathbf{L}_{ij} = \frac{1}{\mathbf{x}_{(i,j)}}$
- Edge incidence matrix $\mathbf{B} \in \mathbf{R}^{|\mathcal{E}| \times |\mathcal{V}|}$, $\mathbf{B}_{(i,j),i} = 1$, $\mathbf{B}_{(i,j),j} = -1$ and $\mathbf{B}_{(i,j),k} = 0$ if $k \neq i, j$
- DC Power Flow Equations become:

$$\mathbf{f} = \mathbf{B} \mathbf{L}^{\dagger} \mathbf{p}$$

Constraints

$$\begin{split} |\mathbf{B}\mathbf{L}^{\dagger}\mathbf{p}| &\leq \overline{\mathbf{f}} \quad (\text{Transmission Capacities}) \\ 0 &\leq \mathbf{p}_{\mathcal{G}_g} \leq \overline{\mathbf{p}}_{\mathcal{G}_g} \quad (\text{Generation Capacities}) \\ &\sum_{i} \mathbf{p}_i = 0 \quad (\text{Power Balance}) \end{split}$$

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- Assume $\mathbf{p}_{\mathcal{G}_r}, \mathbf{p}_{\mathcal{G}_l}$ are fixed to known values
- Set $p_{\mathcal{G}_g}$ so as to minimize "cost" of generation



• **p**⁰: Generation at every node as decided by DCOPF. If **p**_{*G*_{*l*}, **p**_{*G*_{*r*} are fixed in time, DCOPF does the job.}}

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- **BUT**: $\mathbf{p}_{\mathcal{G}_r}$ fluctuates over time: $\mathbf{p}_{\mathcal{G}_r}(t) = \mathbf{p}_{\mathcal{G}_r}^0 + \mathbf{p}_{\mathcal{G}_r}^r(t)$

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- **BUT**: $\mathbf{p}_{\mathcal{G}_r}$ fluctuates over time: $\mathbf{p}_{\mathcal{G}_r}(t) = \mathbf{p}_{\mathcal{G}_r}^0 + \mathbf{p}_{\mathcal{G}_r}^r(t)$
- Controllable generation responds by changing response proportionally

$$\mathbf{p}_{\mathcal{G}_g} = \mathbf{p}_{\mathcal{G}_g}^0 + \alpha_{\mathcal{G}_g} \left(-\sum_{i \in \mathcal{G}_r} \mathbf{p}_i^r(t) \right)$$

Proportions α decided in accordance with generator capacities.

- Add energy storage (batteries) at some nodes $\mathcal{S} \subset \mathcal{V}$
- Can draw power from/supply power to energy storage
- \mathbf{p}_i^s : Power drawn from energy storage at node i ($\mathbf{p}_i^s = 0$ if $i \notin S$)
- \mathbf{s}_i : Energy stored at node i ($\mathbf{s}_i = 0$ if $i \notin S$)



Constraints on Energy Storage

 $0 \leq \mathbf{s}_{\mathcal{S}}$ (Energy cannot be negative)

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• If energy storage is present and being dispatched, proportional control response is:

$$\mathbf{p}_{\mathcal{G}_g} = \mathbf{p}_{\mathcal{G}_g}^0 + \alpha_{\mathcal{G}_g} \left(-\sum_{i \in \mathcal{G}_r} \mathbf{p}_i^r(t) - \sum_{i \in \mathcal{S}} \mathbf{p}_i^s(t) \right)$$

Net Result

$$\mathbf{p}(t) = \mathbf{p}^0 + \mathbf{p}^r(t) + \mathbf{p}^s(t) + \alpha \left(-\sum_{i \in \mathcal{G}_r} \mathbf{p}^r_i(t) - \sum_{i \in \mathcal{S}} \mathbf{p}^s_i(t) \right)$$

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Optimal Control Recipe

- Encode Control Task as Cost Function q(x, u)
- Write down (discrete-time) dynamics $x_{t+1} = f(x_t, u_t)$
- Write down optimization problem:

$$\min_{\substack{x_0, u_0, x_1, u_1, \cdots, x_{T-1}, u_{T-1}, x_T \\ \text{Subject to } x_{t+1} = f(x_t, u_t)} \sum_{t=0}^{T-1} q(x_t, u_t) + q_f(x_T)$$

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Optimize!

Optimal Control Example: Pendulum Swing Up



$$\begin{aligned} x &= (\theta, \dot{\theta}), u = \ddot{\theta} \\ \frac{\theta_{t+1} - \theta_t}{\Delta} &= \dot{\theta}_t \\ \frac{\dot{\theta}_{t+1} - \dot{\theta}_t}{\Delta} &= u - g \sin(\theta) \\ q &= (1 + \cos(\theta))^2 + u^2 \end{aligned}$$

Optimal Energy Storage Control: Cost Function



 cli (p): Penalize transmission lines overloading: cli (p) = ∑_{(i,j)∈E} f(M^{ij}p, -f̄_{ij}, f̄_{ij})
 cg (p^{ns}): Penalize controllable generators overloading/underloading: cg (p^{ns}) = ∑_{i∈G_g} f(p^{ns}_i, 0, p̄_i)
 csp (p^s): Penalize storage use:

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$$csp(\mathbf{p}^s)$$
: Penalize storage use
 $csp(\mathbf{p}^s) = \sum_{i \in S} h(\mathbf{p}^s_i)$

Optimal Energy Storage Control with Perfect Forecasts

Compute optimal energy storage dispatch $\mathbf{p}^{s*}(t)$ by solving

$$\begin{split} \min_{\mathbf{p}_{\mathcal{S}}^{s}(0),...,\mathbf{p}_{\mathcal{S}}^{s}(\mathcal{T}_{f}-1)} &\sum_{t=0}^{\mathcal{T}_{f}} \operatorname{cli}\left(\mathbf{p}\right) + \operatorname{cg}\left(\mathbf{p}^{ns}\right) + \operatorname{csp}\left(\mathbf{p}^{s}\right) \\ \text{subject to} \\ \mathbf{s}(t) &= \sum_{\tau=0}^{t-1} \mathbf{p}^{s}(t) \Delta, \mathbf{s}(\mathcal{T}_{f}) = \mathbf{s}(0), \mathbf{p}_{i}^{s}(t) = 0 \quad i \notin \mathcal{S} \\ \mathbf{p}(t) &= \mathbf{p}^{0} + \mathbf{p}^{s}(t) + \mathbf{p}^{r}(t) + \alpha \left(-\sum_{i \in \mathcal{G}_{r}} \mathbf{p}_{i}^{r}(t) - \sum_{i} \mathbf{p}_{i}^{s}(t)\right) \\ \mathbf{p}^{ns}(t) &= \mathbf{p}(t) - \mathbf{p}^{s}(t) \end{split}$$

- Can be recast as unconstrained problem using $\mathbf{p}_{\mathcal{S}}^{s}(t+1) = \frac{\mathbf{s}_{\mathcal{S}}(t+1)-\mathbf{s}_{\mathcal{S}}(t)}{\Delta}$
- Leads to sparse block-structured Hessian: Efficient Newton's method
- Use Levenberg-Marquardt correction to achieve convergence in practice

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Choose thresholds \epsilon, \epsilon'
\mathcal{S} \leftarrow \mathcal{V}
repeat
    for k = 1 \rightarrow N do
        Generate Random Time Series Profiles for the Renewables
        Solve optimal control problem for the given profiles to get optimal
        dispatch \mathbf{p}^{s*}(t)
        A^k \leftarrow \max_t |\mathbf{p}^{s*}(t)|
    end for
   A \leftarrow \frac{1}{N} \sum_{k=1}^{N} A^k
    \gamma \leftarrow \max\{\gamma : \{\operatorname{perf}(\{i \in \mathcal{S} : A_i \ge \gamma \max(A)\}) < \operatorname{perf}(\mathcal{S}) + \epsilon\}\}
   \mathcal{S} \leftarrow \{i \in \mathcal{S} : A_i > \gamma \max(A)\}
until \gamma < \epsilon'
```

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A more principled approach

$$\begin{split} \min_{\mathbf{p}^{\bar{s}}, \bar{\mathbf{s}}} \frac{1}{N} \sum_{i=1}^{N} \operatorname{Opt}(\bar{\mathbf{p}^{s}}, \bar{\mathbf{s}}, \mathbf{p}^{ri}) + \operatorname{SparsitySizingPenalty}(\bar{\mathbf{p}^{s}}, \bar{\mathbf{s}}) \\ \operatorname{Opt}(\bar{\mathbf{p}^{s}}, \bar{\mathbf{s}}, \mathbf{p}^{r}) &= \min_{\mathbf{p}^{s}(0), \dots, \mathbf{p}^{s}(T_{f}-1)} \sum_{t=0}^{T_{f}} \operatorname{cli}(\mathbf{p}) + \operatorname{cg}(\mathbf{p}^{ns}) + \operatorname{csp}(\mathbf{p}^{s}) \\ \operatorname{Subject to} \\ \mathbf{s}(t) &= \sum_{\tau=0}^{t-1} \mathbf{p}^{s}(t)\Delta, 0 \leq \mathbf{s} \leq \bar{\mathbf{s}}, \mathbf{s}(T_{f}) = \mathbf{s}(0), |\mathbf{p}^{s}| \leq \bar{\mathbf{p}^{s}} \\ \mathbf{p}(t) &= \mathbf{p}^{0} + \mathbf{p}^{s}(t) + \mathbf{p}^{r}(t) + \alpha \left(-\sum_{i \in \mathcal{G}_{r}} \mathbf{p}_{i}^{r}(t) - \sum_{i} \mathbf{p}_{i}^{s}(t)\right) \\ \mathbf{p}^{ns}(t) &= \mathbf{p}(t) - \mathbf{p}^{s}(t) \end{split}$$

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Snapshot of Test Grid

Grid Structure



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Square nodes selected, Red:Active, Green:Inactive

Iteration 1

Activity



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Renewable Penetration

$$\frac{\sum_{t=0}^{T_f} \sum_{i \in \mathcal{G}_r} \mathbf{p}_i^{ns}(t)}{\sum_{t=0}^{T_f} \sum_{i \in \mathcal{G}_l} \mathbf{p}_i^{ns}(t)}$$

Normalized Power Capacity

$$\frac{\sum_{j \in \mathcal{S}} \max_t \left| \mathbf{p}^{s*}_j(t) \right|}{\sum_{i \in \mathcal{G}_r} (\max_t \mathbf{p}^r_i(t) - \min_t \mathbf{p}^r_i(t))}$$

Normalized Energy Capacity

$$\frac{\sum_{j \in \mathcal{S}} (\max_t \mathbf{s}_j^*(t) - \min_t \mathbf{s}_j^*(t))}{\sum_{i \in \mathcal{G}_r} (\max_t \mathbf{s}_i^r(t) - \min_t sr_i(t))} \quad \text{wh}$$

where
$$\mathbf{s}_i^r(t) = \int_0^T \mathbf{p}_i^r(t) dt$$

- 4 cases considered:
 - Unconstrained: Storage at All Nodes
 - Constrained: Storage at 10 nodes (iteration 2 of algorithm)
 - Highly Constrained: Storage at 2 nodes (iteration 3 of algorithm)
 - Intuitive Placement: Storage at Renewables

Experimental Setup

2000 trials for each case, each trial consists of:

- Generate a random penetration value p between 0 and .5
- Generate a random zero-mean fluctuation profile for each renewable:

$$\mathbf{p}_i^r(t) = \sum_k rac{1}{k} \sin\left(\omega_k(t\Delta + \tau) + \phi_k
ight)$$

where ω_k are harmonics of a base frequency and ϕ_k are randomly generated phases.



• Scale the fluctuations by the mean and add: $\mathbf{p}_i^{ns}(t) = \mathbf{p}_i^0(t) * (1 + \mathbf{p}_i^r(t)), i \in \mathcal{G}_r$

Solve the Optimal control problem

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A LOT!

Constraint Violation vs Penetration



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Smart Grid Student Seminar 2

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Energy Capacity vs Penetration

Power Capacity vs Penetration

Average Constraint Violation vs Penetration







Smart Grid Student Seminar

General Observations

- Control necessary to maintain grid constraits under fluctuations
- Energy Storage Control solves the problem upto about 30% penetration
- Even a little fluctuation needs control, proportional control fails to maintain all constraints
- Energy/Power capacity required remain constant upto 30%, then shoot up
- Conjecture: Energy/Power capacity requirements can be pushed down with "smarter" control of traditional generators

Placement Heuristic

- Placing storage at renewables works
- However, requires higher energy/power capacity
- Placement heuristic seems to discover "optimal" placement
- Placed on lines connecting "North" and "South" of grid

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- Control necessary to maintain grid stability at High Penetration Levels
- Operations-driven Planning necessary for Cost-Efficient Renewable Integration
- Optimal Control a useful paradigm in answering this question
- Our initial results indicate that successful Renewable Integration upto 30% Penetration is possible using optimal energy storage control

- Investigate principled approach to placement heuristic
- Test algorithm on various graphs, relative placement of renewables, numbers of renewable nodes
- Combine energy storage control with "smart" generation control
- Build feedback controllers to deal with imperfect predictions
- Decentralized/Distributed solutions?

Questions?

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