

PONDERING ABOUT

QPCP's

(QUANTUM PROBABILISTICALLY
CHECKABLE PROOFS)

OR:

CAN ENERGY GAPS
OF LOCAL HAMILTONIANS
BE AMPLIFIED?

DORET AHARONOV

HEBREW UNIVERSITY

ONGOING WORK WITH

ITAI ARAD

U.C. BERKELEY

PCP

$$|E| = m$$



NP-HARD TO DECIDE:

$$\frac{g_V}{m} = 0 \quad \text{OR} \quad \frac{g_V}{m} \geq \frac{1}{m}$$

PCP THEOREM



G



G'

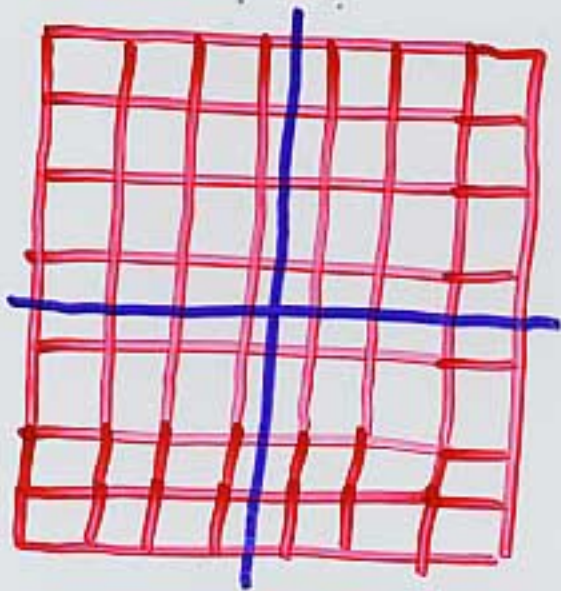
$$\frac{g_V(G)}{m} = 0 \quad \rightarrow \quad \frac{g_V(G')}{m'} = 0$$

$$\frac{g_V(G)}{m} \geq \frac{1}{m} \quad \rightarrow \quad \frac{g_V(G')}{m'} = \text{const!}$$

NP
HARD
!

INDEED, PROBABILISTICALLY
CHECKABLE PROOF!

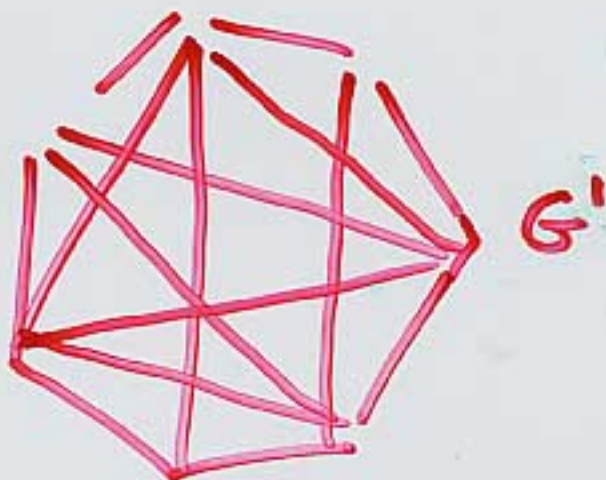
BUT...



ISN'T THIS
A POLYTIME
ALG FOR

$g \pm \frac{1}{\epsilon} \cdot m$?

NO PCP IN DIM = CONST



PCP WORKS WITH HIGHLY
CONNECTED GRAPHS (EXPANDERS..)

HISTORY

BABAI, FORTNOW, LEVEN, SZEGEDY } [91, 92]
BABAI, FORTNOW, LUND }
GOLDWASSER, MICALI, RACKOFF }

ARORA, SAFRA }
ARORA, LUND } $q = \text{CONST!}$
MOTWANI, SUDAN, SZEGEDY }

BELLARE, GOLDWASSER, LUND, RUSSELL
BELLARE, COPPERSMITH, HASTAD, KJWI, SUDAN

9↓

HASTAD $q = 3 \text{ BITS!}$

NEW PROF: DINUR [06]

QUANTUM NP [KITAEV 93]



$$H = \sum_{i \sim j} H_{ij} \quad \leftarrow \text{PROJECTIONS}$$

LOCAL HAMILTONIAN PROBLEM

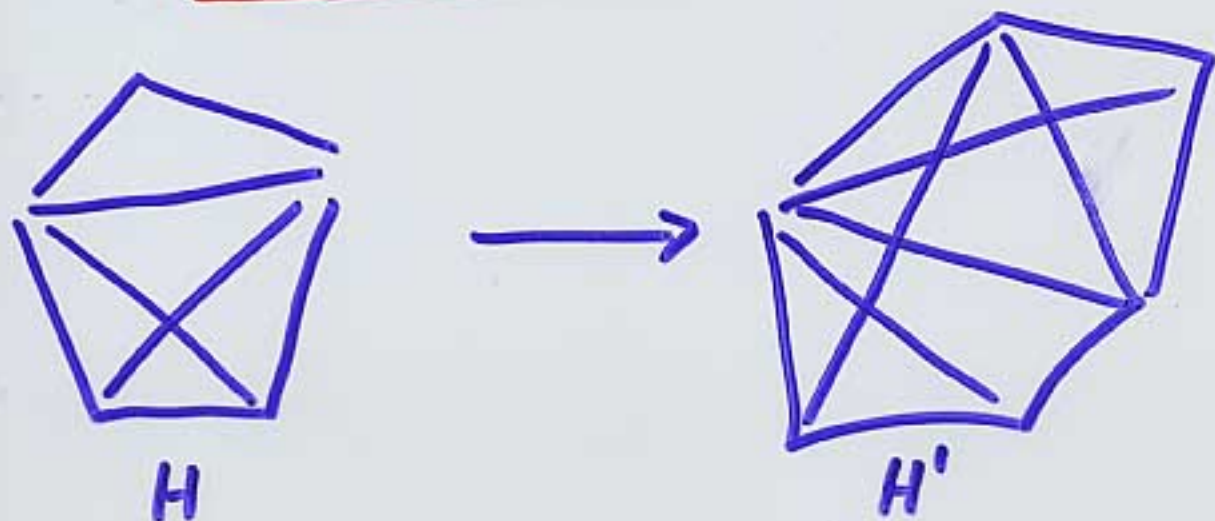
$$g\psi(H) = 0 \quad \text{OR} \quad g\psi(H) \geq \frac{1}{\text{poly}(m)}$$

CAN BE VERIFIED EASILY (BY A QUANTUM VERIFIER!)

IF CORRECT: JUST GIVE THE $|g\rangle$.

LOCAL HAMILTONIAN IS QNP-COMPLETE.

QUANTUM PCP ?



$$g\psi(H) = 0$$

$$g\psi(H) \geq \frac{1}{m}$$

$$\longrightarrow g\psi_{m'}(H') = 0$$

$$\longrightarrow g\psi_{m'}(H') \geq \text{CONST!}$$

POSSIBLE IMPLICATIONS OF QPCP

- 1) Inability to Approx $g_s(H) \pm \epsilon_0 \mu$
(EVEN WITH A Q. COMPUTER)
- 2) SPECTRAL GAP AMPLIFICATIONS.

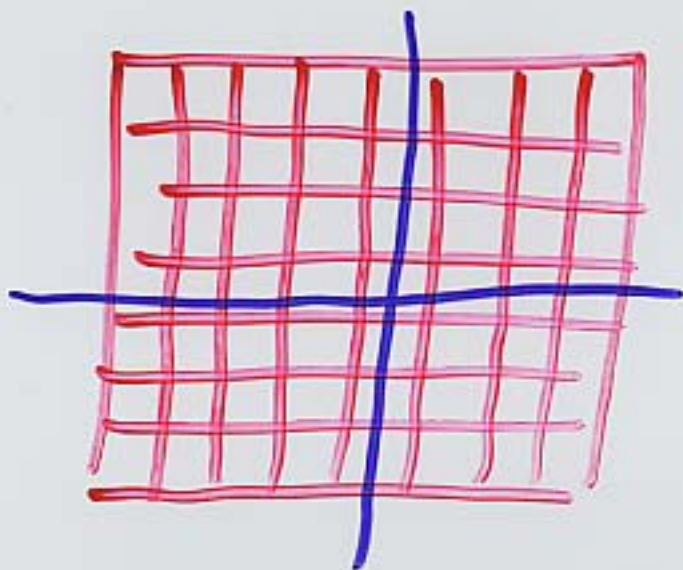
PROMISE GAP \neq SPECTRAL GAP.
IN QPCP

BUT... DINUR'S CONSTRUCTION
CAN BE MODIFIED TO AMPLIFY
SPECTRAL GAP.

* $H \rightarrow H'$ W. AMPLIFIED
SPECTRAL GAP.

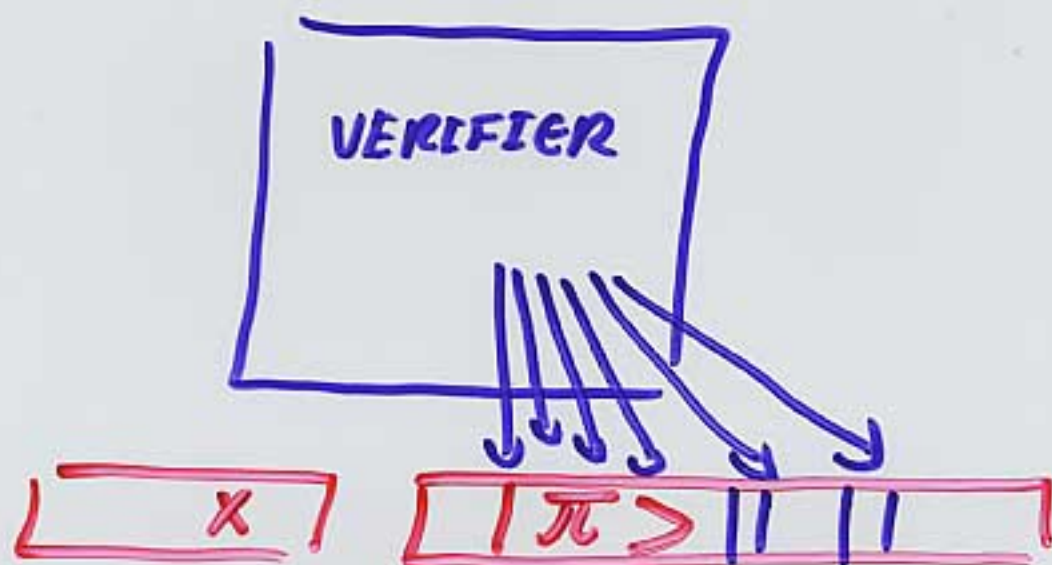
* IMPLICATIONS TO FAULT TOLERANT
ADIABATIC COMPUTATION.

ONCE AGAIN - DIM \neq CONST



IF DIM BOUNDED - CAN ALWAYS
APPROX TO WITHIN $\pm \epsilon_0 m \forall \epsilon_0$.

MAYBE WE CAN PROVE WEAKER QPCP THEOREMS?



1) MAYBE ALLOW V TO ACCESS $\log n$
LOCATIONS?

2) MAYBE PROVE

$$QMA = \left\{ g^U = 0 \text{ OR } g^U > \frac{1}{\log n} \right\}$$

3) MAYBE ALLOW V ACCESS TO
RANDOMLY CHOSEN $g = \text{CONST}$ QUBITS,
BUT GIVE HIM THE FULL MATRIX?
(MIXING-QPCP)

OUTLINE OF REST OF TALK:

DINUR'S PROOF -

A SKETCH ...

OUR RESULTS.

QPCP DIFFICULTIES

ONE OF THE MORE INTERESTING
IMPLICATIONS - SETTLING QPCP

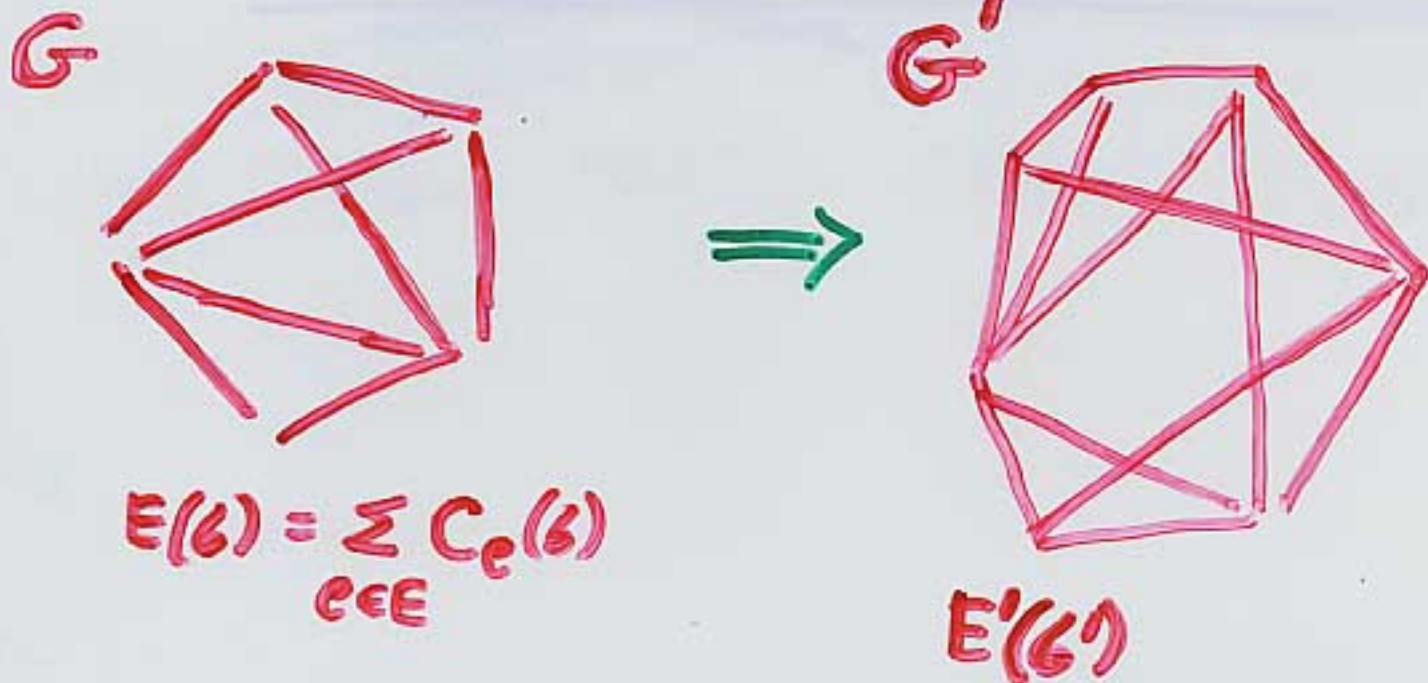
BETTER UNDERSTANDING
OF ENTANGLEMENT
& GROUND STATES.

SEEMS TO CLASH WITH NO-CLONING
BUT NO PROOF OF THAT!

BEEN IN THAT STORY..

13 YEARS AGO, QECC SEEMED TO
DO THAT TOO...

GENERAL OUTLINE



s.t:

$$g_{\text{v}}(G) = 0 \Rightarrow g_{\text{v}}(G') = 0$$
$$g_{\text{v}}(G) > 0 \Rightarrow g_{\text{v}}(G') \geq \epsilon \cdot m'$$

CENTRAL NOTION: EXPANDERS



DEGREE (VALENCY) = d (CONST)
YET - BEHAVES RANDOMLY.
MIXES RAPIDLY...

ON ONE HAND - MIXES ERRORS. DISTRIBUTES...
ON THE OTHER - $|E| = d \cdot |V|$

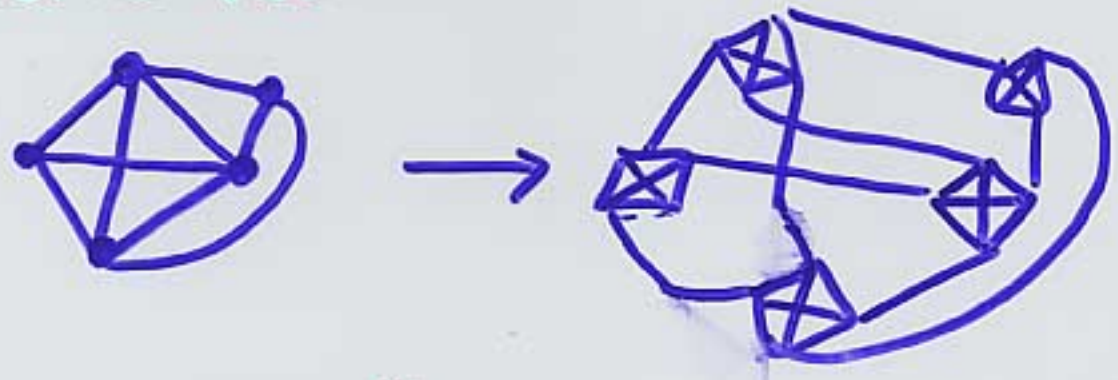
USE EXPANDERS FOR $\log(m)$ ITERATIONS.

$$\frac{1}{m} \xrightarrow{\times 2} \dots \xrightarrow{\times 2} \text{CONST.} \quad \& \quad \text{SIZE}' = \text{POLY}'$$

ONE ITERATION:

$$g \rightarrow g \times \text{CONST}$$
$$|G| \rightarrow |G| \times \text{CONST}$$

1) DEGREE REDUCTION



2) "EXPANDERIZE": $G \rightarrow GU$ EXPANDER

CONSISTENCY CONSTRAINTS \Rightarrow EXPANDER OF SMALL DEGREE, $g \sim$ O.K

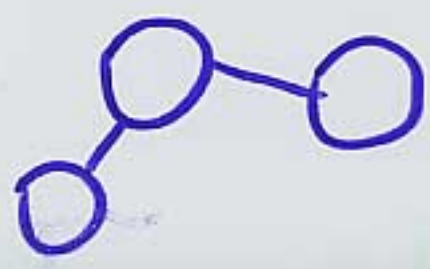
3) GAP AMPLIFICATION.

EXPANDER
SMALL
DEGREE



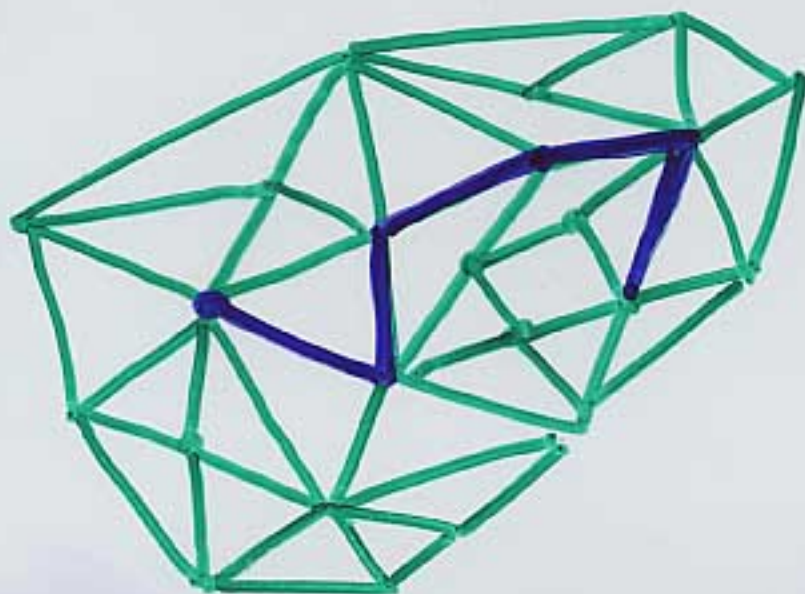
$G', g \sim(G') > g \sim(G)$
BUT: $|\Sigma|$ HUGE!

4) ALPHABET REDUCTION.



GAP AMPLIFICATION

MAIN IDEA: t -WALKS ON EXPANDERS,



t -steps.

EACH t -WALK \Rightarrow NEW CONSTRAINT.

RAPID MIXING \Rightarrow MANY t -WALKS WILL
"CATCH" THE PROBLEM.
Factor of t amplification!

BUT - NO LONGER A GRAPH!

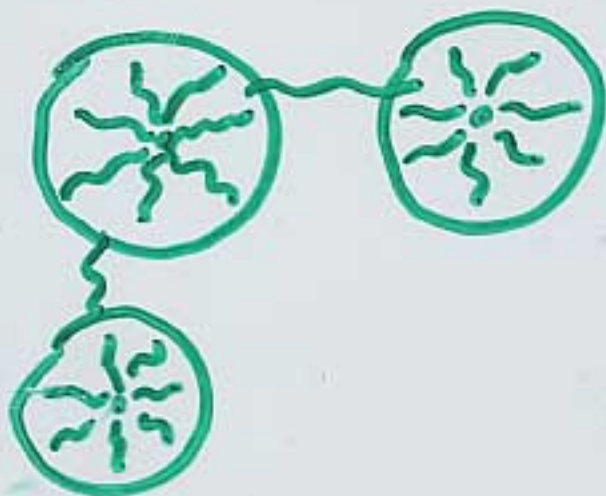
NEW VERTICES:

SPHERES OF t -WALKS,
 $|S| = d^t$

NEW CONSTRAINT:

CONSISTENCY,

& INHERITED CONSTRAINTS



CAN WE GENERALIZE TO \mathcal{Q} ?

Main Problem: Lots of consistency checks!

SEEMS TO VIOLATE NO-CLONING THM...

ISOLATE OTHER ESSENTIAL INGREDIENT:



t -WALKS \rightarrow



CONSTRAINT:
 $N \gg E \gg B \gg C \dots$

$$H_{t\text{-walk}} = I - P_{A \rightarrow B \rightarrow C \dots}$$

$$H' = \sum_{t\text{-walks}} H_{t\text{-walk}}$$

HOW MUCH LARGER IS $\frac{g \nu(H')}{m^2}$

COMPARED TO $\frac{g \nu(H)}{m}$?

HMMM...

EXAMPLE $t=2$.



$$A = |00\rangle\langle 00|, \quad B = |00\rangle\langle 00|.$$

1) $|\psi\rangle = |001\rangle.$

$$\frac{E}{m} = \frac{1}{2}.$$

$$\frac{E'}{m'} = \frac{1}{1} = 1.$$



2) $|\psi\rangle = \sqrt{1-\epsilon}|000\rangle + \sqrt{\epsilon}|111\rangle.$

$$\frac{E}{m} = \frac{\epsilon + \epsilon}{2} = \epsilon.$$

$$\frac{E'}{m'} = \epsilon!$$



CAPYTALYST CASE - OK

SOCIALIST CASE - NO AMPLIFICATION!

QUANTUM (WEAK) GAP AMPLIFICATION

FOR t -WALKS:

$$g_{\psi}^{(H')} > \left[1 + o\left(\frac{t}{\log(m)}\right) \right] \frac{g_{\psi}^{(H)}}{m}$$

MAIN OBSERVATION - SOCIALIST CASE CAN'T HAPPEN.



$$|\psi\rangle = |\alpha\rangle + |\beta\rangle + |\gamma\rangle$$

$\underbrace{\hspace{1cm}}_{A \cap B} \quad \underbrace{\hspace{1cm}}_{(A \cup B)^c} \quad \underbrace{\hspace{1cm}}_{\text{REST}}$

NO AMPLIFICATION:

$$\mathcal{E} = E_A(|\psi\rangle) = \|\beta\rangle\|^2 + \|(I - P_A)|\gamma\rangle\|^2$$

$$\mathcal{E} = E_B(|\psi\rangle) = \|\beta\rangle\|^2 + \|(I - P_B)|\gamma\rangle\|^2$$

$$\mathcal{E} = E_{A \cap B}(|\psi\rangle) = \|\beta\rangle\|^2 + \|\gamma\rangle\|^2$$

$$\Rightarrow |\gamma\rangle = 0!$$

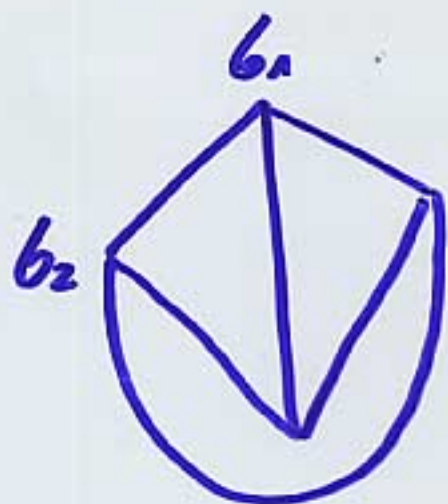
$$\Rightarrow P_A |\psi\rangle \in A \cap B$$

SIMILARLY: $P_A |\psi\rangle \in A \cap C!$

$$\Rightarrow P_A |\psi\rangle \in A \cap B \cap C \dots \Rightarrow g_{\psi} = 0!$$



NP



$$E(b_1, b_2, \dots, b_n) = \sum_{i \sim j} C_{ij}(b_i, b_j)$$

$$\exists \vec{b} \text{ s.t. } E(\vec{b}) = 0 ?$$

CAN EASILY VERIFY $E(\vec{b}) = 0$ GIVEN \vec{b} .

NP PROBLEMS: EASY TO VERIFY GIVEN
A HINT (MAYBE HARD TO
FIND!)

ANOTHER EXAMPLE:

CORRECT MATHEMATICAL THEOREMS

& THOUSAND OTHER EXAMPLES...

NP-COMPLETENESS.

e.g: GRAPH COLORABILITY IS NP-COMPLE

$$X \rightarrow G$$

$$\left. \begin{array}{l} X \in \text{YES} \rightarrow g_v(G) = 0 \\ X \in \text{NO} \rightarrow g_v(G) \geq 1. \end{array} \right\} \text{HARD PROBLEM !}$$

CONCLUSIONS

- QPCP IS AN IMPORTANT & FUNDAMENTAL QUESTION.
(IN BOTH DIRECTIONS)
- HOPE YOU HAVE GOTTEN SOME FLAVOR OF DINUR'S BEAUTIFUL PROOF!
- INTERESTING PRELIMINARY QUANTUM RESULTS.